

**Higher Engineering Mathematics**  
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**Lecture - 50**  
**Duality - I**

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**Duality Concept**

One of the most interesting concepts in linear programming is the duality theory. Every linear programming problem has associated with it, another linear programming problem involving the same data and closely related optimal solutions. Such two problems are said to be duals of each other. While one of these is called the primal, the other the dual.

The importance of the duality concept is due to two main reasons. Firstly, if the primal contains a large number of constraints and a smaller number of variables, the labour of computation can be considerably reduced by converting it into the dual problem and then solving it. Secondly, the interpretation of the dual variables from the cost or economic point of view proves extremely useful in making future decisions in the activities being programmed.



Hello friends, welcome to my lecture on The Concept of Duality, one of the most important and interesting concepts in linear programming is the Duality. Every linear programming problem has associated with it, another linear programming problem which involves the same data and is closely related to the optimal solutions, such two problems are said to be duals of each other.

While one of these problems is called the primal problem, the other one is called the dual problem. The importance of the duality concept is due to the two main reasons. Firstly, if the primal contains a large number of constraints and a smaller number of variables then the labour of computation can be considerably reduced by converting it into the dual problem and then solving it, because if there are say  $m$  constraints and  $n$  variables and  $m$  is greater than  $n$  then when we will consider the corresponding dual problem there the number of constraints will become  $n$  and number of variables will become  $m$ , ok.

So, number of constraints will reduce when we consider the dual problem and therefore it will be easier to solve the get the solution of the dual problem. So labour of computation can be reduced by converting the problem into the corresponding dual problem and then solving

it. Secondly, the interpretation of the dual variables from the cost or economic point of view proves extremely useful in making future decisions in the activities being programmed.


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**Formation of the dual problem**

Consider the following L.P.P. :

**Maximize**  $Z = c_1x_1 + c_2x_2 + \dots + c_nx_n$ ,  
 subject to the constraints

$a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n \leq b_1$ ,  
 $a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n \leq b_2$ ,  
 .....  
 $a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n \leq b_m$ ,  
 $x_1, x_2, \dots, x_n \geq 0$ .



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Now, let us consider the following L.P.P. maximize  $Z = c_1 x_1 + c_2 x_2$  and so on  $c_n x_n$  subject to these  $m$  constraints  $a_{11} x_1 + a_{12} x_2$  and so on  $a_{1n} x_n \leq b_1$ ;  $a_{21} x_1 + a_{22} x_2$  and so on  $a_{2n} x_n \leq b_2$  and then the  $m$  at constraint  $a_{m1} x_1 + a_{m2} x_2$  and so on  $a_{mn} x_n \leq b_m$ ,  $x_1, x_2, x_n$  are all non-negative.

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To construct the dual problem, we adopt the following guidelines:

- (i) The maximization problem in the primal becomes the minimization problem in the dual and vice versa.
- (ii) ( $\leq$ ) type of constraints in the primal becomes ( $\geq$ ) type of constraints in the dual and vice versa.
- (iii) The coefficients  $c_1, c_2, \dots, c_n$  in the objective function of the primal become  $b_1, b_2, \dots, b_m$  in the objective function of the dual.
- (iv) The constraints  $b_1, b_2, \dots, b_m$  in the constraints of the primal become  $c_1, c_2, \dots, c_n$  in the constraints of the dual.
- (v) If the primal has  $n$  variables and  $m$  constraints, the dual will have  $m$  variables and  $n$  constraints i.e. the transpose of the body matrix of the primal problem gives the body matrix of the dual.
- (vi) The variables in both the primal and dual are non-negative.


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### Formation of the dual problem

Consider the following L.P.P. :

$$\text{Maximize } Z = c_1x_1 + c_2x_2 + \dots + c_nx_n,$$

subject to the constraints

$$a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n \leq b_1,$$

$$a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n \leq b_2,$$

$$\dots$$

$$a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n \leq b_m,$$

$$x_1, x_2, \dots, x_n \geq 0.$$

$$\begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \dots & \dots & \dots & \dots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{bmatrix}_{m \times n}$$

$$\begin{bmatrix} a_{11} & a_{21} & \dots & a_{m1} \\ a_{12} & a_{22} & \dots & a_{m2} \\ \dots & \dots & \dots & \dots \\ a_{1n} & a_{2n} & \dots & a_{mn} \end{bmatrix}_{n \times m}$$



Now, in order to construct the dual problem we adopt the following guidelines, the maximization problem in the primal becomes the minimization problem in the dual and vice versa.  $\leq$  type of constraints in the primal become  $\geq$  type of constraints in the dual and vice versa.

The coefficients  $c_1, c_2, \dots, c_n$  in the objective function, the coefficients  $c_1, c_2, \dots, c_n$  in the objective function of the primal become  $b_1, b_2, \dots, b_m$ , ok so  $c_1, c_2, \dots, c_n$  will be replaced by  $b_1, b_2, \dots, b_m$ , ok these  $b_1, b_2, \dots, b_m$  and  $x_1, x_2, \dots, x_n$  will be replaced by new variables  $y_1, y_2, \dots, y_m$ , ok.

So, the dual problem in the dual problem maximization will be replaced by minimization and we will have minimization of  $Z^i$  where we will have  $Z^i = b_1 y_1 + b_2 y_2$  and so on  $b_m y_m$ , ok.

If the primal has  $n$  variables and  $m$  constants like here we have  $n$  variables  $x_1, x_2, \dots, x_n$  and  $m$  constraints given by these  $m$  an equalities, ok so then the dual will have  $m$  variables and  $n$  constraints that is the transpose of the body matrix of the primal problem gives the body matrix of the dual.

So, here the body matrix is this one, body matrix is  $a_{11}, a_{12}$  and so on  $a_{1n}, a_{21}, a_{22}, \dots, a_{2n}$  and so on  $a_{m1}, a_{m2}$  and so on  $a_{mn}$  this is body matrix, ok of the primal problem.

In the dual problem this will be replaced with  $a_{11}, a_{21}$  and so on  $a_{m1}, a_{12}, a_{21}, a_{m2}$ , ok and then lastly  $a_{1n}, a_{2n}$  and so on  $a_{mn}$ , ok so here we have here we have  $m$  rows and  $n$  columns, so this

is m by n matrix, here we have n rows and we can see we have m there are 1, 2 and so on m columns and n rows, so n by m, ok.

So, here the body matrix, it is m by n size, here it is of the size n by m, transpose of the body matrix of the primal problem, ok. So transpose of the body matrix of the primal problem gives the body matrix of the dual.

The variables in both the primal and dual are non-negative, ok. So, now the dual problem will then be you see here, the dual problem will now be  $Z^i$  minimization of  $Z^i = b_1 y_1 + b_2 y_2$  and so on  $b_m y_m$ , subject to the constraints  $a_{11} y_1 + a_{21} y_2$  and so on  $+ a_{m1} y_m \geq c_1$ , ok  $a_{12} y_1, a_{22} y_2 \dots + a_{m2} y_m \geq c_2$  and so on  $a_{1n} y_1 + a_{2n} y_2$  and so on  $a_{mn} y_m \geq c_n$ , ok.

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Then, the dual problem will be

**Minimize**  $Z = b_1 y_1 + b_2 y_2 + \dots + b_m y_m$   
 subject to the constraints

$a_{11} y_1 + a_{21} y_2 + \dots + a_{m1} y_m \geq c_1,$   
 $a_{12} y_1 + a_{22} y_2 + \dots + a_{m2} y_m \geq c_2,$   
 .....  
 $a_{1n} y_1 + a_{2n} y_2 + \dots + a_{mn} y_m \geq c_n,$   
 $y_1, y_2, \dots, y_m \geq 0.$

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**Formation of the dual problem**

Consider the following L.P.P. :

**Maximize**  $Z = c_1 x_1 + c_2 x_2 + \dots + c_n x_n,$   
 subject to the constraints

$a_{11} x_1 + a_{12} x_2 + \dots + a_{1n} x_n \leq b_1,$   
 $a_{21} x_1 + a_{22} x_2 + \dots + a_{2n} x_n \leq b_2,$   
 .....  
 $a_{m1} x_1 + a_{m2} x_2 + \dots + a_{mn} x_n \leq b_m,$   
 $x_1, x_2, \dots, x_n \geq 0.$

$$\begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & & a_{2n} \\ \dots & \dots & \dots & \dots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{bmatrix}_{m \times n}$$

$$\begin{bmatrix} a_{11} & a_{21} & \dots & a_{m1} \\ a_{12} & a_{22} & & a_{m2} \\ \dots & \dots & \dots & \dots \\ a_{1n} & a_{2n} & \dots & a_{mn} \end{bmatrix}_{n \times m}$$

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So, we will have these equations a 1 Z minimization of  $Z = b_1 y_1, b_2 y_2, b_m y_m$  subject to the constants  $a_{11} y_1 + a_{21} y_2, \dots + a_{mn} y_m \geq c_n$  will be replaced with  $\geq$  ok in the dual. So, we will have  $\geq$  here, so  $\geq c_1, \geq c_2$  and  $\geq c_n$  and  $y_1, y_2, y_m$  are non-negative variables, ok.  $\leq$

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**Example**

Let us find the dual of the following L.P.P.

**Minimize**  $Z = 3x_1 - 2x_2 + 4x_3$ ,  
 subject to  $3x_1 + 5x_2 + 4x_3 \geq 7$ ,  
 $6x_1 + x_2 + 3x_3 \geq 4$ ,  
 $7x_1 - 2x_2 - x_3 \leq 10$ ,  
 $x_1 - 2x_2 + 5x_3 \geq 3$ ,  
 $4x_1 + 7x_2 - 2x_3 \geq 2$ ,  
 $x_1, x_2, x_3 \geq 0$ .

*Since the problem is of minimization, all the constraints must be  $\geq$  type.*

*Maximize  $W = 7y_1 + 4y_2 - 10y_3 + 3y_4$*   
*Subject to*  
 $3y_1 + 6y_2 - 7y_3 + y_4 \leq 3$   
 $5y_1 + y_2 + 2y_3 - 2y_4 \leq -2$   
 $4y_1 + 3y_2 + y_3 + 5y_4 \leq 4$   
 $y_1, y_2, y_3 \geq 0$

*body matrix =*

3	5	4
6	1	3
-7	2	1
1	-2	5

$\Rightarrow$  Transpose =

3	6	-7	1
5	1	2	-2
4	3	1	5

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Let us find the dual of the following L.P.P. So, we have  $Z = 3x_1 - 2x_2 + 4x_3$  subject to  $3x_1 + 5x_2 + 4x_3 \geq 7$ , since the problem is of minimization you can see here, since the problem is of minimization all the constraints must be of  $\geq$  type, all the constraints must be  $\geq$  type, ok

so  $\geq$  type.

So here what is happening in this constraint, this is  $\leq$  type, so we have to convert it to  $\geq$  type, so when you convert it to  $\geq$  type it will become  $-7x_1 + 2x_2 + x_3 \geq -10$ , ok  $-7x_1 + 2x_2 + x_3 \geq -10$ , ok.

Now, so let us write the dual problem, so maximization of maximize Z will be replaced by some W, W = now this is  $c_1, c_2, c_3$  W replaced by with  $b_1, b_2, \dots b_m$ , so  $7y_1 + 4y_2 - 10y_3$ , ok and then  $3y_4$ , ok and then we will have so this is our objective function in the dual problem and (corres subs) subject to the constraints, we will have now the (bod) the matrix here is body this body matrix is 3, 5, 4, ok 6, 1, 3, ok then we have -7, 2, 1, ok we have 1 - 2, 5 this is body matrix in the primal problem.

So, in the dual problem we have to take transpose of this, so we will have subject to constants  $3y_1 + 6y_2 - 7y_3 + y_4$  and  $\leq$  we will have  $\leq c_1, c_1$  is 3, ok then we have  $5y_1 + y_2 + 2y_3$  and we

get  $-2y_4 \leq -2$ , ok and then  $4y_1 + 3y_2 + y_3 + 5y_4 \leq 4$ , ok so  $c_1, c_2, c_3$ , ok and  $y_1, y_2, y_3 \geq 0$ , so this is our dual problem, we replaced (mak) minimization by maximization, the variable  $Z$  will be replaced by some other variable,

some other variable, so we can write it as  $W$ .

Then the  $c_1, c_2, c_3$  are replaced with  $b_1, b_2, b_3, b_4$ , ok so  $b_1$  is 7,  $b_2$  is 4,  $b_3$  is -10,  $b_4$  is 3, ok and we use new variables  $y_1, y_2, y_3, y_4$ , so  $7y_1 + 4y_2 - 10y_3 + 3y_4$ , the body matrix here is 3, 5, 4, 6, 1, 3 - 7, 2, 1 and then 1 - 2, 5 we take transpose of this and then it becomes transpose of this becomes 3, 6, -7, 1, 5, 1, 2, -2, 4, 3, 1, 5, ok.

So, we get the constraints as  $3y_1 + 6y_2 - 7y_3 + y_4 \leq c_1$ ,  $5y_1 + y_2 + 2y_3 - 2y_4 \leq -2$ ,  $4y_1 + 3y_2 + y_3 + 5y_4 \leq 4$ , so this and where  $y_1, y_2, y_3$  are non-negative, so this is the dual problem, ok.

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minimize  $W = b_1 y_1 + b_2 y_2$   
 Subject to  $a_{11} y_1 + a_{21} y_2 \geq c_1$   
 $a_{12} y_1 + a_{22} y_2 \geq c_2$   
 where  $y_1, y_2 \geq 0$   
 $b_1, b_2$  is unrestricted  
 $c_1, c_2$  is nonnegative

**Formulation of the dual problem when the primal has equality constraints**

Consider the problem  
**Maximize**  $Z = c_1 x_1 + c_2 x_2$   
 subject to  $a_{11} x_1 + a_{12} x_2 = b_1$   
 $a_{21} x_1 + a_{22} x_2 \leq b_2$   
 $x_1, x_2 \geq 0$ .

The equality constraints can be written as  
 $a_{11} x_1 + a_{12} x_2 \leq b_1$  and  $a_{11} x_1 + a_{12} x_2 \geq b_1$ ,  
 or  $a_{11} x_1 + a_{12} x_2 \leq b_1$  and  $-a_{11} x_1 - a_{12} x_2 \leq -b_1$ ,

$$\begin{bmatrix} a_{11} & a_{12} \\ -a_{11} & -a_{12} \\ a_{21} & a_{22} \end{bmatrix}^T = \begin{bmatrix} a_{11} & -a_{11} & a_{21} \\ a_{12} & -a_{12} & a_{22} \end{bmatrix}$$

minimize  $W = b_1 y_1' - b_1 y_1'' + b_2 y_2$   
 subject to  
 $a_{11} y_1' - a_{11} y_1'' + a_{21} y_2 \geq c_1$   
 $a_{12} y_1' - a_{12} y_1'' + a_{22} y_2 \geq c_2$   
 where  $y_1', y_1'', y_2 \geq 0$   
 Let us define  $y_1 = y_1' - y_1''$   
 then  $y_1$  is unrestricted in sign

Now, let us consider the problem where we have to maximize  $Z = c_1 x_1 + c_2 x_2$  and what we have here the change here is that 1 constraint is an equality, ok there is a equality constraints  $a_{11} x_1 + a_{12} x_2 = b_1$ , so how we will (siles) find the dual of this as L.P.P. ok. So, what we will do is the equality constraint  $a_{11} x_1 + a_{12} x_2 = b_1$  can be written as  $a_{11} x_1 + a_{12} x_2 \leq b_1$  and  $a_{11} x_1 + a_{12} x_2 \geq b_1$ , so in the problem is of maximization, ok all the constraints should be of  $\leq$  type, ok.

So, first constraint will be  $a_{11} x_1 + a_{12} x_2 \leq b_1$ , the second constraint this one will be converted into  $\leq$  type and we will write it as  $-a_{11} x_1 - a_{12} x_2 \leq -b_1$ . So, now the problems will this be replaced by  $a_{11} x_1 + a_{12} x_2 \leq b_1$ ,  $-a_{11} x_1 - a_{12} x_2 \leq -b_1$ , ok and  $a_{21} x_1, a_{22} x_2 \leq b_2$ , ok.

Now, so when we write the dual problem here we write maximize maximization we replaced by minimization minimize  $W =$  now here we have  $b_1, -b_1, b_2$ , ok so we write  $b_1$ , ok we can write  $b_1 y_1'$ , ok  $b_1 c_1, c_2$  will be replaced with  $b_1, -b_1$  and  $b_2$ , so  $b_1 y_1' - b_1 y_1''$ , let us take the variables  $y_1', y_1''$  and then  $b_2 y_2$ , ok.

So, this and the in the body matrix is now  $a_{11}, a_{12}$ , ok  $-a_{11}, -a_{12}, a_{21}, a_{22}$  this is  $3 \times 3$  matrix, ok 3 rows and columns, so we take the transpose of this body matrix, so transpose of this will be giving you  $a_{11}, -a_{11}, a_{21}$  and then  $a_{12}, -a_{12}$  and then  $a_{22}$ , ok so subject to constraints  $a_{11} y_1' - a_{11} y_1'' + a_{21} y_2 \geq c_1$  we will have because we have  $\leq$  here, so  $\geq c_1$ , ok then  $a_{12} y_1' - a_{12} y_1'' + a_{22} y_2 \geq c_2$ , so where  $y_1', y_1''$  and  $y_2$  are all  $\geq 0$ .

Now, let us define  $y_1 = y_1' - y_1''$ , since  $y_1'$  and  $y_1''$  are greater than both non-negative  $y_1$  is unrestricted in sign, ok it is not restricted in sign, so  $y_1$  then  $y_1$  is unrestricted in sign and we will get this problem as minimization of, so this will then change to minimization of  $W = b_1 y_1' - y_1''$  will be replaced by  $y_1$ , so  $b_1 y_1 + b_2 y_2$ , ok subject to  $a_{11}$  into  $y_1' - y_1''$ .

So,  $a_{11} y_1 + a_{21} y_2 \geq c_1$  and then  $a_{12}$  multiplied by  $y_1' - y_1''$  that is  $y_1$ , so  $a_{12} y_1 + a_{22} y_2 \geq c_2$  where  $y_2 \geq 0$  and  $y_1$  is unrestricted in sign, so  $y_1$  can be negative as well as positive, so this is the dual of the given problem when the primal has equality constraints we will do this, ok.

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Ctd...

Then the above problem can be restated as

**Maximize**  $Z = c_1 x_1 + c_2 x_2$   
 subject to  $a_{11} x_1 + a_{12} x_2 \leq b_1$  and  $-a_{11} x_1 - a_{12} x_2 \leq -b_1$ ,  
 $a_{21} x_1 + a_{22} x_2 \leq b_2$ ,  
 $x_1, x_2 \geq 0$ .

Now we form the dual using  $y_1', y_1'', y_2$  as the dual variables. Then the dual problem is

**Minimize**  $W = b_1(y_1' - y_1'') + b_2 y_2$  ✓  
 subject to  $a_{11}(y_1' - y_1'') + a_{21} y_2 \geq c_1$ ,  
 $a_{12}(y_1' - y_1'') + a_{22} y_2 \geq c_2$ ,  
 $y_1', y_1'', y_2 \geq 0$ . ✓


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So, this what we have then, this is maximize  $Z =$  this, subject to this, ok this constraints  $x_1, x_2 \geq 0$  then we use the variables  $y_1', y_1'', y_2$ , ok and we get  $W$ ,  $W =$  this, subject to this, ok and which gives us this problem, ok.



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Ctd....

The term  $(y_1' - y_1'')$  appears in both the objective function and all the constraints of the dual. This will always happen whenever there is an equality constraint in the primal. Then the new variable  $(y_1' - y_1'')$  becomes unrestricted in sign being the difference of two non-negative variables and the above dual problem takes the form.

**Minimize**  $W = b_1 y_1 + b_2 y_2,$  ✓  
 subject to  $a_{11} y_1 + a_{21} y_2 \geq c_1,$  ✓  
 $a_{12} y_1 + a_{22} y_2 \geq c_2, y_1$  unrestricted in sign,  $y_2 \geq 0.$

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The term  $y_1' - y_1''$  appears both in the objective function and all the constraints of the dual. This will always happen whenever there is an equality constraint in the primal. The new variable this becomes unrestricted in sign, ok because of the difference of two non-negative variables, so the above dual problem then takes this form, ok which we have just now seen, ok.

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Ctd....

In general, if the primal problem is

**Maximize**  $Z = c_1 x_1 + c_2 x_2 + \dots + c_n x_n,$   
 subject to the constraints

$a_{11} x_1 + a_{12} x_2 + \dots + a_{1n} x_n = b_1,$   
 $a_{21} x_1 + a_{22} x_2 + \dots + a_{2n} x_n = b_2,$   
 .....  
 $a_{m1} x_1 + a_{m2} x_2 + \dots + a_{mn} x_n = b_m,$   
 $x_1, x_2, \dots, x_n \geq 0.$

Minimize  $W = b_1 y_1 + b_2 y_2 + \dots + b_m y_m$   
 Subject to  
 $a_{11} y_1 + a_{21} y_2 + \dots + a_{m1} y_m \geq c_1$   
 $a_{12} y_1 + a_{22} y_2 + \dots + a_{m2} y_m \geq c_2$   
 .....  
 $a_{1n} y_1 + a_{2n} y_2 + \dots + a_{mn} y_m \geq c_n$   
 where  $y_1, y_2, \dots, y_m$   
 are unrestricted in sign

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Now, let us consider in general if the primal problem is  $Z = c_1 x_1 + c_2 x_2$  and so on  $c_n x_n$  subject to this constraints, this is the  $= b_1$ , this is the  $= b_2$ , this is the  $= b_m$ , ok then it will be

replaced with the it is dual of this will be minimize  $W$ ,  $W = b_1 y_1 + b_2 y_2$  and so on  $b_m y_m$ , ok  $b_m y_m$  subject to the constraints  $a_{11}$ , yeah  $a_{11} y_1 + a_{21} y_2$  and so on  $a_{m1} y_m$ , ok.

We will have because it is maximization problems, so we will have  $\leq$  here we will have  $\geq c_1$ , ok then  $a_{12} y_1 + a_{22} y_2$  and so on  $a_{m2} y_m \geq c_2$  and so on, we shall have  $a_{1n} y_1, a_{2n} y_2$  and so on  $a_{mn} y_m \geq c_m$ , ok  $c_n$  sorry  $c_n$  where  $y_1, y_2, \dots, y_m$  are all unrestricted in sign, so this is the dual, ok.

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Then the, dual problem will be  
**Minimize**  $W = b_1 y_1 + b_2 y_2 + \dots + b_m y_m$   
 subject to the constraints  
 $a_{11} y_1 + a_{21} y_2 + \dots + a_{m1} y_m \geq c_1$   
 $a_{12} y_1 + a_{22} y_2 + \dots + a_{m2} y_m \geq c_2$   
 .....  
 $a_{1n} y_1 + a_{2n} y_2 + \dots + a_{mn} y_m \geq c_n$   
 $y_1, y_2, \dots, y_m$  all unrestricted in sign.  
 Thus the dual variables corresponding to the equality constraints are unrestricted in sign. Conversely, when the primal variables are unrestricted in sign, corresponding dual constraints are equalities.

So, the dual will be  $W = b_1 y_1 + b_2 y_2$  and so on  $b_m y_m$  subject to this constraints, ok where these are all unrestricted in sign, thus the dual variables corresponding to the equality constraints are unrestricted in sign. Conversely, when the primal variables are unrestricted in sign corresponding to dual constraints are equalities, ok so if so when the when we have equality constraints the dual variables are unrestricted in sign and if the primal variables unrestricted in sign then the corresponding dual constraints are equalities, ok.

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$$\begin{bmatrix} 2 & 3 & 2 \\ 3 & -2 & 4 \\ -3 & 2 & -4 \end{bmatrix}$$



**Example:**

Construct the dual of the L.P.P.:

**Maximize**  $Z = 4x_1 + 9x_2 + 2x_3$ ,  
 subject to  $2x_1 + 3x_2 + 2x_3 \leq 7$ ,  
 $3x_1 - 2x_2 + 4x_3 = 5$ ,  
 $x_1, x_2, x_3 \geq 0$ .

*Handwritten notes:*  
 $3x_1 - 2x_2 + 4x_3 \leq 5$   
 $3x_1 - 2x_2 + 4x_3 \geq 5$  or  $-3x_1 + 2x_2 - 4x_3 \leq -5$

*Minimize*  $W = 7y_1 + 5y_2 - 5y_2'$  or  $7y_1 + 5y_2$   
 Subject to  
 $2y_1 + 3y_2 - 3y_2' \geq 4$   
 $3y_1 - 2y_2 + 2y_2' \geq 9$  where  $y_1, y_2 \geq 0$   
 $2y_1 + 4y_2 - 4y_2' \geq 2$  where  $y_2'$  is unrestricted in sign and  $y_2 \geq 0$



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Now, let us construct the dual of the L.P.P.  $Z = (mik) 4x_1 + 9x_2 + 2x_3$  subject to  $2x_1 + 3x_2 + 2x_3 \leq 7$  and  $3x_1 - 2x_2$ , so there is a 1 equality constraints here, so we have here we can write it as  $3x_1 - 2x_2 + 4x_3 \leq 5$ , ok and  $3x_1 - 2x_2 + 4x_3 \geq 5$ , we will have to convert to  $\leq$  because the problem is of maximization, so  $-3x_1 + 2x_2 - 4x_3 \leq -5$ , ok.

So, we have the three constraints one is this one, ok  $2x_1 + 3x_2 + 2x_3 \leq 7$  and other one  $3x_1 - 2x_2 + 4x_3 \leq 5$  and third one is  $-3x_1 + 2x_2 - 4x_3 \leq -5$ , so  $\leq -5$ , so the dual problem will be minimization of  $W = 7y_1$ , ok  $7y_1 + 5y_2' - 5y_2''$ , ok and then we shall have the body matrix is what here? 2, 3, 2, ok 3, -2, 4, ok and then we have -3, 2, -4, so this is body matrix, ok we have to take the corresponding transpose for the body matrix of the dual problems.

So, we get 2, 3, -3, so  $2y_1 + 3y_2' - 3y_2''$ , ok and we will have to put  $\geq 4$ , ok and then we have  $3y_1 - 2y_2$ , so this should be written as  $y_2$  and  $y_2''$  not y sorry, this should be, ok so, ok. So,  $3y_1 - 2y_2' + 2y_2'' \geq 9$  and then we have  $2y_1 + 4y_2' - 4y_2'' \geq 2$ , ok or we can write it as where  $y_1, y_2', y_2''$  they are all non-negative, ok.

Now, we can write the this problem also in the form minimize  $W = 7y_1 + 5y_2$ , we will write  $y_2 = y_2' - y_2''$ , ok and then subject to  $2y_1 + 3y_2 \geq 4$ ,  $3y_1 - 2y_2 \geq 9$  and  $2y_1 + 4y_2 \geq 2$  where  $y_2 = y_2' - y_2''$ , ok and so  $y_2$  is unrestricted in sign, ok.

So, the dual problem is minimize  $W = 7y_1 + 5y_2$  subject to the constraints  $2y_1 + 3y_2 \geq 4$ ,  $3y_1 - 2y_2 \geq 9$ ,  $2y_1 + 4y_2 \geq 2$ ,  $y_2$  is unrestricted in sign,  $y_1 \geq 0$ . So, this is how we write the dual of this L.P.P. where 1 constraint is of equality type. So that is all in this lecture, thank you very much for your attention.