

**Higher Engineering Mathematics**  
**Prof. P.N. Agrawal**  
**Department of Mathematics**  
**Indian Institute of Technology Roorkee**  
**Various types of Graphs - I**  
**Mod06\_Lec26**

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**Introduction**

Many situations that occur in Computer Science, Physical Science, Communication Science, Economics and many other areas can be analysed by using techniques found in a relatively new area of mathematics called graph theory.

Hello friends, welcome to my lecture on various types of Graphs. Let us first see why we need graph theory, many situations that occur in computer science, physical science, communication science, economics and many other areas can be analysed by using techniques found in a relatively new area of mathematics called graph theory.

(Refer Slide Time: 0:53)

The graph can be used to represent almost any problem involving discrete arrangements of objects, where concern is not with the internal properties of these objects but with relationship among them. We begin with basic graph terminology and then will discuss some important concepts in graph theory with many applications of graphs.

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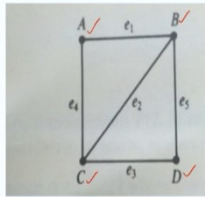
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### Graph

**Definition** A graph  $G$  consists of two parts:

- (i) A set  $V = V(G)$  whose elements are called vertices, points or nodes.
  - (ii) A collection  $E = E(G)$  of unordered pairs of distinct vertices called edges.
- We write  $G(V, E)$  when we want to emphasize the two parts of  $G$ . Consider the graph shown in the figure below.

### Example



In the above graph  $G = G(V, E)$ ,  $V$  consists of four vertices A, B, C, D and  $E$  consist of five edges  $e_1 = \{A, B\}$ ,  $e_2 = \{B, C\}$ ,  $e_3 = \{C, D\}$ ,  $e_4 = \{A, C\}$ ,  $e_5 = \{B, D\}$ .



Now, let us first define what we mean by a graph, a graph consist of two parts  $V = V(G)$  okay, whose elements are called vertices, points or nodes okay, a graph is  $G$ , this is graph is  $G$ . A graph  $G$  consist of two parts, a set  $V = V(G)$  whose elements are called vertices, points or nodes. A collection  $E = E(G)$  of unordered pairs of vertices called edges. We write  $G = G(V, E)$  when we want to emphasise the two parts of  $G$ . Let us consider the graph shown in this figure, okay.

Here you can see that, there are four vertices in this is graph. Okay, A, B, C and D. Okay and A, B, C, D are joined by means of the edges  $e_1$ ,  $e_1$  is  $\{A, B\}$ ,  $e_4$  that is  $\{A, C\}$ ,  $e_2$  that is  $\{B, C\}$ ,  $e_3$  that is  $\{C, D\}$  and  $\{B, D\}$ , that is  $e_5$  okay, so there are four vertices A, B, C, D and E. Okay, consist E is the set of edges, E consist of five edges  $e_1, e_2, e_3, e_4, e_5$ , so this is a graph.

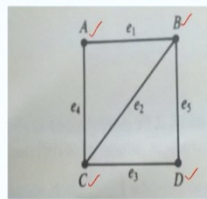
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### Simple graph

A graph that has neither self-loops nor parallel edges is called a simple graph.



### Example



*Simple graph*

In the above graph  $G = G(V, E)$ ,  $V$  consists of four vertices  $A, B, C, D$  and  $E$  consist of five edges  $e_1 = \{A, B\}$ ,  $e_2 = \{B, C\}$ ,  $e_3 = \{C, D\}$ ,  $e_4 = \{A, C\}$ ,  $e_5 = \{B, D\}$ .



Now, let us see what is a simple graph? A graph that has neither self-loops nor parallel edges okay are called a simple graph. So this graph is a simple graph, this graph is a simple graph, because it does not contain parallel edges okay, now are self-loops, what are simple parallel ages and loops.

(Refer Slide Time: 3:08)

### Multigraph

In a multigraph, each vertex  $V$  is represented by a dot (or a small circle) and each edge  $E = \{u, v\}$  is represented by a curve which connects end points  $u$  and  $v$ . (In fact, we usually denote a graph, when possible, by drawing its diagram rather than explicitly listing its vertices and edges.)



Let us see, we first define multigraph, in a multigraph, each vertex  $V$  is represented by a dot or a small circle and each edge  $E = \{u, v\}$  is represented by a curve which connects end points  $u$  and  $v$ , a graph is generally shown by drawing its diagram rather than by explicitly listing its vertices and edges.

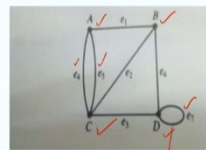
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**Parallel edges in a graph:** All edges having the same pair of end vertices are called parallel edges.

**Self loop in a graph:** An edge having the same vertices is called a self loop.



### Example



- $e_1 = (A, B)$
- $e_2 = (B, C)$
- $e_3 = (C, D)$
- $e_4 = (A, C)$
- $e_5 = (A, C)$
- $e_6 = (B, D)$
- $e_7 = (D, D)$

Self loop  
 $e_4$  and  $e_5$  parallel edges



So, we have talked about parallel edges, what are parallel edges? All edges having the same pair of end vertices are called parallel edges. And self-loop is an edge having the same vertices is called a self-loop. Let us look at this figure, in this is a multigraph, you can see here there are four vertices A, B, C and D. Okay and the edges are  $e_1$ . Okay,  $e_1$ , which is (A, B) and then we have  $e_2$ ,  $e_2$  is (B, C) and then  $e_3$ ,  $e_3$  is (C, D),  $e_4$ ,  $e_4$  is (A, C),  $e_5$ ,  $e_5$  is (A, C) again,  $e_6$  (B, D),  $e_7$  is (D, D).

So, you can see that this edge  $e_7$  has same end vertices okay, D, D same end vertices, so this is a self-loop and  $e_4$  and  $e_5$  edges have the same end vertices A, C. Okay,  $e_4$  and  $e_5$  have the

same end vertices, so  $e_4$  and  $e_5$  are parallel edges, so this graph consist of parallel edges,  $e_4$  and  $e_5$  and self-loop  $e_7$  okay.

(Refer Slide Time: 5:36)

**Example**

Consider the diagram of the multigraph  $G(V, E)$  where  $V = \{P_1, P_2, P_3, P_4, P_5\}$  and

$E = \{(P_1, P_1), (P_2, P_3), (P_2, P_4), (P_3, P_2), (P_4, P_1), (P_5, P_4)\}$

*(P<sub>2</sub>, P<sub>3</sub>) & (P<sub>3</sub>, P<sub>2</sub>) are parallel edges*

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10

Now, let us go to this diagram okay, here again we see that there are 1, 2, 3, 4, 5, 5 vertices okay, vertices are  $P_1, P_2, P_3, P_4, P_5$  okay and the set of edges is E equal to  $(P_1, P_1), (P_1, P_1)$  is a because  $(P_1, P_1)$  means it has the same end vertices okay, so same, but this is here, so this is a self-loop, so this is a self-loop because this edge has same end vertices  $(P_1, P_1)$  okay, now we have  $(P_2, P_3)$ . Okay,  $(P_2, P_3)$  and we have  $(P_3, P_2)$  okay, since here again the end vertices are same, these are parallel edges okay, so  $(P_2, P_3)$  and  $(P_3, P_2)$  are parallel edges okay and we have  $(P_2, P_4)$ ,  $(P_2, P_4)$ , we have  $(P_4, P_1)$  okay,  $(P_4, P_1)$  here,  $(P_5, P_4)$  okay, so this is a multigraph which consist of one self-loop and two parallel edges.

(Refer Slide Time: 7:03)

### Finite multigraph

A multigraph  $G = G(V, E)$  is finite if both  $V$  is finite and  $E$  is finite. Note that a graph  $G$  with a finite number of vertices  $V$  must automatically have a finite number of edges and so must be finite.

A multigraph is said to be finite if both the set of vertices is finite and the set of edges is finite, now let us know that if there are only finite number of vertices in the graph that it must automatically having a finite number of edges because when the vertices are finite, when you join them you will have only finite number of edges, so the graph must be finite, so multigraph will be called finite if number of vertices are actually finite.

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**Trivial graph:** A trivial graph is the graph with one vertex and no edges.

**Empty or null graph:** The empty graph is the graph with no vertices and no edges.

**Isolated vertex:** A vertex having no edge incident on it is called an isolated vertex.

Now trivial graph, a trivial graph is the graph with one vertex okay and no edges, so in the case of a trivial graph there is one vertex, no edges. The empty graph is the graph with no vertices and no edges okay. Isolated vertex, a vertex having no edge incident on it is called an isolated vertex.

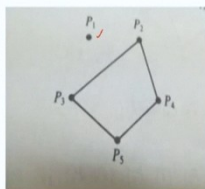


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**Example:** Consider the diagram of the multigraph  $G(V, E)$  where

$V = \{P_1, P_2, P_3, P_4, P_5\}$  and

$E = \{(P_2, P_4), (P_2, P_3), (P_3, P_5), (P_5, P_4)\}$



*P<sub>1</sub> - isolated vertex*

Let us look at this figure. Okay, this graph, so you see here we have five vertices okay,  $P_1, P_2, P_3, P_4, P_5$  okay, five vertices are there, edges are  $(P_2, P_4)$  okay,  $(P_2, P_3)$ ,  $(P_3, P_5)$ ,  $(P_5, P_4)$ , and this  $P_1$ , it has no edges incident on it, so this  $P_1$ , is an isolated vertex because though there are no edges incident on it, so it is an isolated vertex.

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**Example:** Let  $G = G(V, E)$  have five vertices. Find the maximum number  $m$  of edges in  $E$  if:

(a)  $G$  is a graph, and (b)  $G$  is a multigraph.

*no. of edges =  $\binom{5}{2} = 10$*

*simple graph*

Now, let  $G = G(V, E)$  have five vertices okay, let  $G = G(V, E)$  have five vertices. Find the maximum number  $m$  of the edges in  $E$  if  $G$  is a graph okay, since to draw an edge you have to

collect two vertices okay, so  $\binom{5}{2}$ , number of edges, if you take any two vertices okay, if you

take any two vertices then you have number of edges will be  $\binom{5}{2}$ , in a graph, in a graph means a simple graph here okay, we have by graph here we means simple graph, where there are no loops and no parallel edges okay, so you take any two vertices to draw an edge okay, we will have number of edges  $\binom{5}{2}, \binom{5}{2}$  means 10.


So when no self-loops and no parallel edges are there and we have five vertices, number of maximum, maximum number of edges could be  $\binom{5}{2}$  okay, that is 10. Now G is a multigraph, if G is a multi-graph, then you can have any number of parallel edges, any number of self-loops okay, so there is no maximum number of edges in E. Okay, in the case of a multigraph.


(Refer Slide Time: 10:18)

**Definition**

Suppose  $e = \{u, v\}$  is an edge in G, i.e.,  $u$  and  $v$  are endpoints of  $e$ . Then the vertex  $u$  is said to be adjacent to the vertex  $v$ , and the edge  $e$  is said to be incident on  $u$  and on  $v$ .

**Degree and parity (even or odd) of a vertex:** The degree of a vertex  $v$  in a graph G, written  $deg(v)$ , is equal to the number of edges which are incident on  $v$  or, in other words, the number of edges which contain  $v$  as an endpoint. The vertex  $v$  is said to be even or odd according as  $deg(v)$  is even or odd.




15

Now suppose  $e$  equal to  $u, v$  is an graph G. Okay, then  $u$  and  $v$  are endpoints of  $e$ , the vertex  $u$  is called adjacent to the vertex  $v$  and the edge  $e$  is said to be incident on  $u$  and  $v$  on  $v$ . Okay, so if you take  $u$  and  $v$  two points okay, the two vertices and this is an edge okay, then  $u$  and  $v$  are called adjacent to each other and  $e$  is said to be incident on  $u$  and  $v$ .

Degree of a vertex okay, the degree of a vertex could be either even or odd and how do we define the degree? The degree of a vertex  $v$  in a graph written as degree  $v$  is equal to the number of edges which are incident on  $v$  or in other words the number of edges which contain  $v$  as an endpoint, the vertex  $v$  is said to be even or odd according as the degree of  $v$  is even or odd okay.

(Refer Slide Time: 11:24)

**Theorem:** The sum of the degrees of the vertices of a graph is equal to twice the number of edges.

**Proof:** Since each edge is counted twice in counting the degree of the vertices of a graph  $G$ . Therefore the sum of degrees of the vertices of a graph is equal to twice the number of edges.



Now let us see, we have, we see this example. Okay, we will consider this theorem first, the sum of degrees of the vertices of a graph is equal to twice the number of edges, this is because each edge is counted twice. Okay and counting the degree of the vertices of a graph  $G$ , so suppose you have a graph  $G$  and we have two vertices okay,  $u$  and  $v$ , this is the edge okay, then while counting the degree of  $u$  okay, this vertex, this edge will be counted while counting the degree of  $v$ , this edge again will be counted, so since each edge is counted twice in counting the degrees, degree of the vertices of a graph  $G$ . Therefore, the sum of degrees of the vertices of a graph is equal to twice the number of edges.

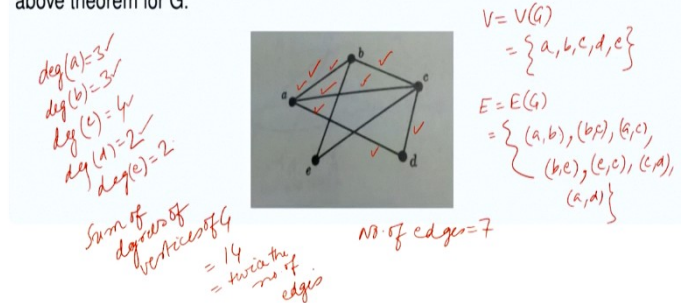
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**Theorem:** The sum of the degrees of the vertices of a graph is equal to twice the number of edges.

**Proof:** Since each edge is counted twice in counting the degree of the vertices of a graph  $G$ . Therefore the sum of degrees of the vertices of a graph is equal to twice the number of edges.



**Example:** Consider the graph  $G = G(V, E)$  in the figure below. (a) Describe  $G$  formally. (b) Find the degree and the parity of each vertex of  $G$ . (c) Verify the above theorem for  $G$ .



Now, let us consider this graph. Okay, this graph  $G = G(V, E)$  in this figure describe  $G$  formally, so this, let us see how many, so in order to describe  $G(V, E)$  we need to know how many vertices are there and how many edges are there, so  $V = V(G)$  this denotes the vertices okay, vertices are  $a, b, c, d, e$ , these are the vertices and then we have  $E = E(G)$  okay.

So we have, now let us list the edges of this multigraph okay, so we have  $a, b$  okay, this is  $a, b$ , then we have  $b, c$ , we have taken this, we have taken this, then we have  $a, c$ , so we have considered this one and then we have  $b, e$ , then we have  $e, c$  and then we have  $c, d$ , we also have  $a, d$  okay, so we have 1, 2, 3, 4, 5, 6 and 7. Okay, seven edges are there okay, now degree of each vertex we have to find okay, so degree of  $a$  let us see, degree of  $a$ , now you see how many edges are incident on  $A$ .

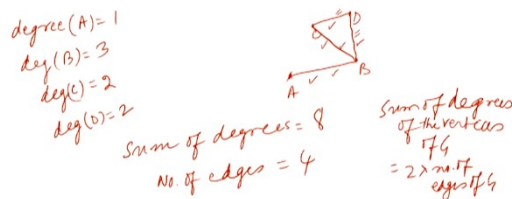
So one that is  $a, b$ , then  $a, c$  and then  $a, d$  3 edges are incident on  $a$ , so degree of  $a$  equal to 3, degree of  $b$ , we have this edge  $b, a$  or you can say  $a, b$  and then we have  $e, b$  and then we have  $c, b$  okay, so 3 edges are incident on  $b$  and then we have degree of  $c$ , so we have  $a, c$ , incident on  $c$ .  $b, c$  and we have  $e, c$ . Okay and we have  $d, c$  also, so we have 4 and then we have degree of let us say  $d$  okay, so we have  $c, d$  incident on  $d$ , we have  $a, d$  incident on  $d$ , so degree of  $d$  is 2 okay and this one and this one.

Now, they are incident on  $d$  and then we have degree of  $e$ , so how many edges are incident on  $e$ .  $b, e$ , and  $c, e$ . Okay, so we have 2 and sum of degrees of the vertices of this multigraph are  $3 + 3 + 6 + 4 = 10, 12, 14$ , sum of degrees = 14 and the edges are 7, number of edges are, so you can see this is twice the number of edges, so this verifies the theorem which we have just

now seen, the sum of the degrees of the vertices of a graph is equal to twice the number of edges.

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**Example:** Consider the graph  $G$  where  $V(G) = \{A, B, C, D\}$  and  $E(G) = \{(A, B), (B, C), (B, D), (C, D)\}$ . Find the degree and parity of each vertex in  $G$ .



Now, let us consider the graph  $G$  where  $V(G) = \{A, B, C, D\}$  so in this graph, there are 4 vertices  $A, B, C$  and we have  $D$ . Okay,  $E(G)$  now this is the set of edges, so we have edges are  $(A, B)$  and we have  $(B, C)$  we have  $(B, D)$  and we have  $(C, D)$  okay, so let us find the degree, so this is a multigraph okay and we have the degree and parity of, degree of each vertex, degree of  $A$ , there is only 1 edge incident on  $A$ , so we have 1, degree of  $B$  equals to, now 1, 2, 3. 3 edges are incident on  $B$ .

So we have 3 degree of  $C$ , there are 2 edges, incident on  $C$ , so we have 2 and degree of  $D$ , there are 1, 2. 2 edges, incident on  $D$ , so we have 2, so total of sum of degrees equal to  $4 + 2 + 2$ , that is 8 okay and edges are 1, 2, 3, 4 okay and number of edges are 4. Okay, so sum of degrees of the vertices of  $G =$  twice the number of edges of  $G$  okay, so that verifies the theorem we have stated. Thank you very much for your attention.