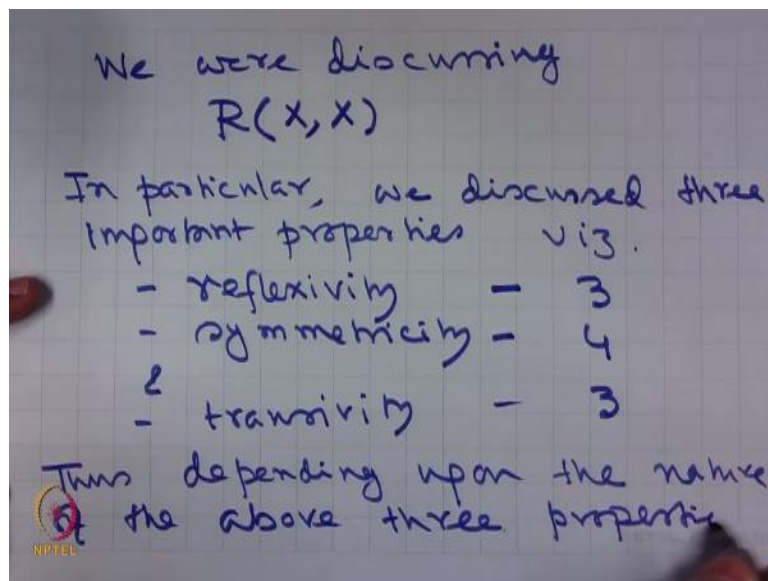


**Introduction to Fuzzy Sets Arithmetic and Logic**  
**Prof. Niladri Chatterjee**  
**Department of Mathematics**  
**Indian Institute of Technology – Delhi**

**Lecture - 20**  
**Fuzzy Sets Arithmetic and Logic**

Welcome students to the MOOCs course on fuzzy sets arithmetic and logic, this is lecture number 20.

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In the last class we were discussing binary relations between the same set that is  $R(X, X)$ .

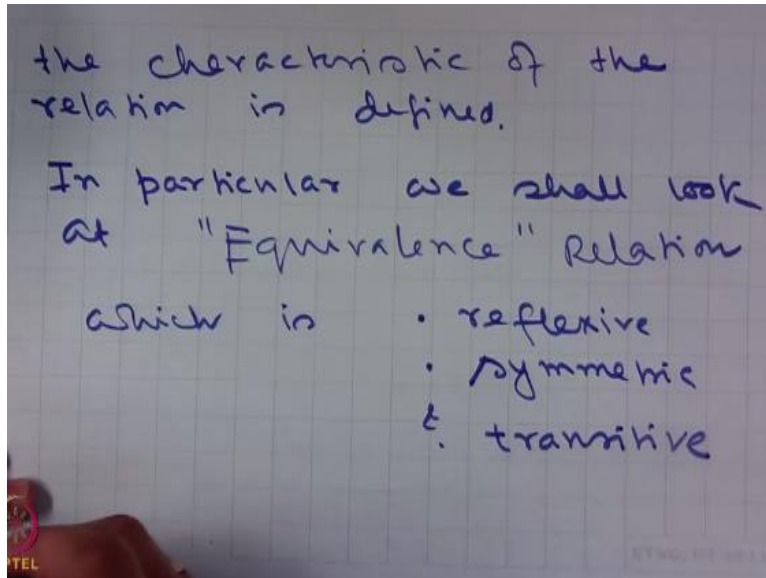
In particular, we discussed three important properties namely

- reflexivity
- symmetry
- transitivity

Reflexivity has 3 variations, Symmetry has 4 variations and transitivity also has 3 variations.

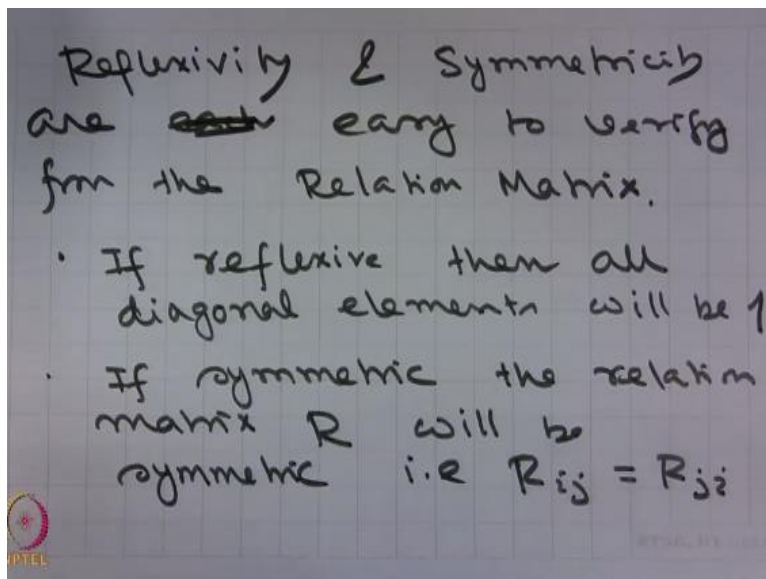
Thus, depending upon the nature of the above three properties the characteristic of the relation is defined.

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In particular, we look at equivalence relation which is reflexive, symmetric and transitive. For some other relations such as partially ordered set or strictly order set you can proceed in a very similar way and see how those relationships can be defined with respect to crisp sets and also fuzzy sets.

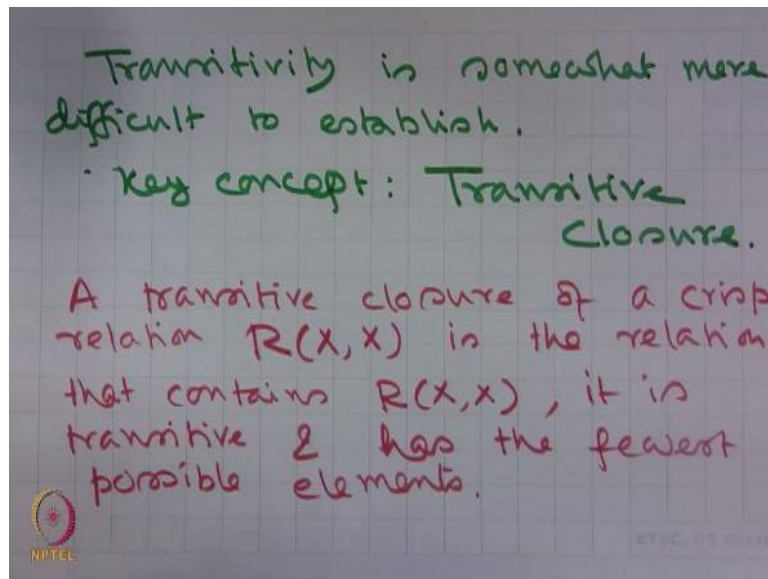
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Now reflexivity and symmetricity are easy to verify from the relation matrix.

- If it is reflexive, then all diagonal elements will be 1 and
- If symmetric the relation matrix  $R$  will be symmetric, that is  $R_{ij} = R_{ji}$  where  $R_{ij}$  is the element at the  $i^{th}$  row and  $j^{th}$  column and  $R_{ji}$  is the element in the  $j^{th}$  row,  $i^{th}$  column.

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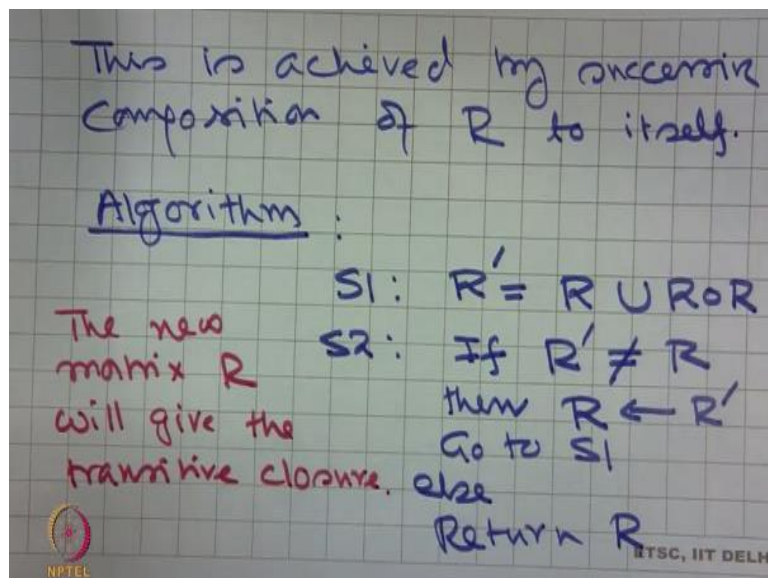
However, transitivity is somewhat more difficult to establish.

The key concept here is transitive closure.

A transitive closure of a crisp relation  $R(X,X)$  is the relation that contains  $R(X,X)$ , it is transitive and has the fewest possible elements,

That is given a relation  $R$  from  $X$  to  $X$ , the transitive closure is the smallest set that contains  $R$  and also that is transitive.

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This is achieved by successive composition of  $R$  to itself.

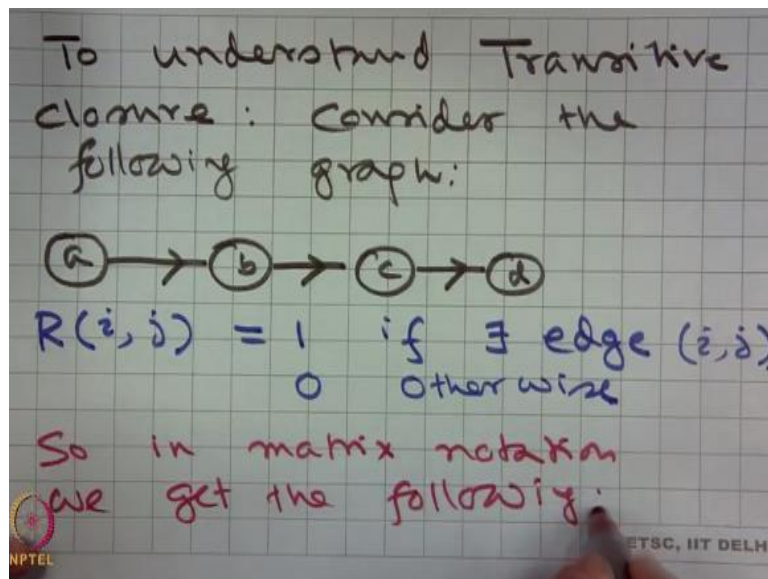
So there is a simple algorithm for that one for achieving transitive closure.

Step 1.  $R' = R \cup (R \circ R)$

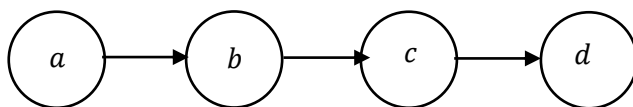
Step 2. If  $R' \neq R$  then,  $R \leftarrow R'$  and go to Step 1,  
else return  $R$ .

The new matrix  $R$  will give the transitive closure.

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To understand the meaning of transitive closure consider the following graph.

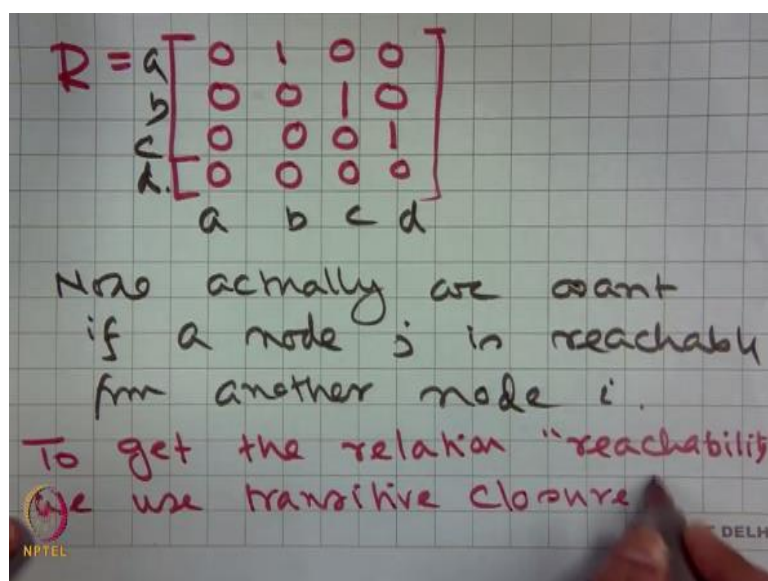


There are four nodes  $a, b, c$  and  $d$ . From  $a$  we can reach  $b$ , from  $b$  we can reach  $c$  and from  $c$  we can reach  $d$ . So,

$$R(i, j) = \begin{cases} 1 & \text{if } \exists \text{ edge}(i, j) \\ 0 & \text{otherwise} \end{cases}$$

So in matrix notation we get the following.

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$$R = \begin{matrix} a & & 0 & 1 & 0 & 0 \\ b & & 0 & 0 & 1 & 0 \\ c & & 0 & 0 & 0 & 1 \\ d & & 0 & 0 & 0 & 0 \end{matrix} \begin{bmatrix} & a & b & c & d \\ & & & & \end{bmatrix}$$

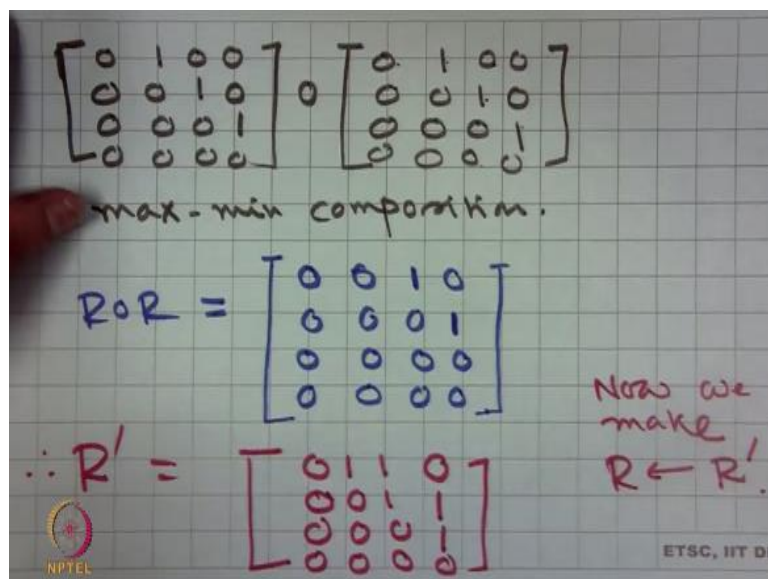
This shows that there is an edge from  $a$  to  $b$ , there is an edge from  $b$  to  $c$  and there is an edge from  $c$  to  $d$ .

Now, actually we want if a node  $j$  is reachable from another node  $i$ .

And if we go back to the graph we see that  $c$  is reachable from  $a$  via  $b$ ,  $d$  is also reachable from  $a$  via  $b$  and  $c$  and  $d$  is also reachable from  $b$ .

Therefore, to get the relation reachability we use transitive closure.

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So we compose  $R$  with itself and we use max-min composition. So we get the following matrix.

$$R \circ R = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \circ \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$\therefore R' = R \cup (R \circ R) = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

Since it is different from original  $R$  now we make  $R \leftarrow R'$

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We now compose  $R \circ R$

$$R = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \cup \begin{bmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$= \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$\therefore R' = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$

We now compose this new  $R$  with itself

$$R \circ R = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \circ \begin{bmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

Therefore, new  $R'$  after the second iteration is

$$\therefore R' = R \cup (R \circ R) = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \cup \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

Since it is different from these  $R$ , now this matrix is going to be the new  $R$  and I make a composition with itself.

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$$\begin{aligned}
 & \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \circ \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \\
 & = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \\
 & = R' = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}
 \end{aligned}$$

∴ The transitive closure is

Therefore,

$$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \circ \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

Again if I take the union new  $R'$  is going to be

$$R' = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

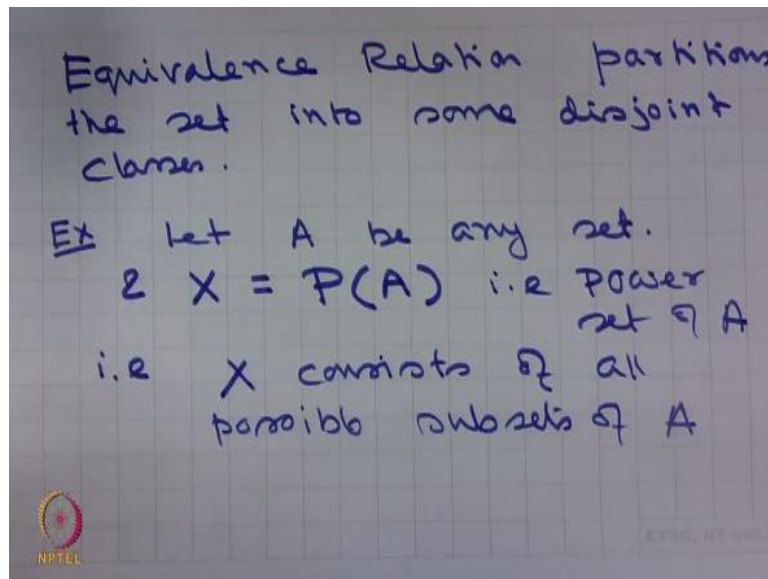
Thus we get this  $R'$  is same as this  $R$ .

Therefore, the transitive closure is this matrix and it shows that

- $b$  is reachable from  $a$ ,
- $c$  is reachable from  $a$ ,
- $d$  is reachable from  $a$ ,
- $c$  is reachable from  $b$ ,
- $d$  is reachable from  $b$  and
- $d$  is reachable from  $c$ .

So that gives the idea of transitive closure.

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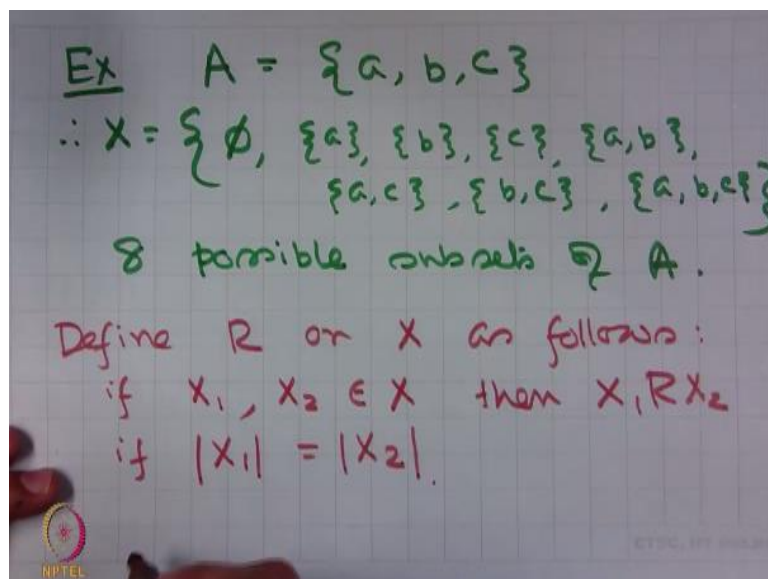


Equivalence relation partitions the set into some disjoint classes.

Let me give you an example.

Let  $A$  be any set and  $X = P(A)$  that is the power set of  $A$ , that means  $X$  consists of all possible subsets of  $A$ .

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Example

Let  $A = \{a, b, c\}$

Therefore,  $X = \{ \emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\} \}$

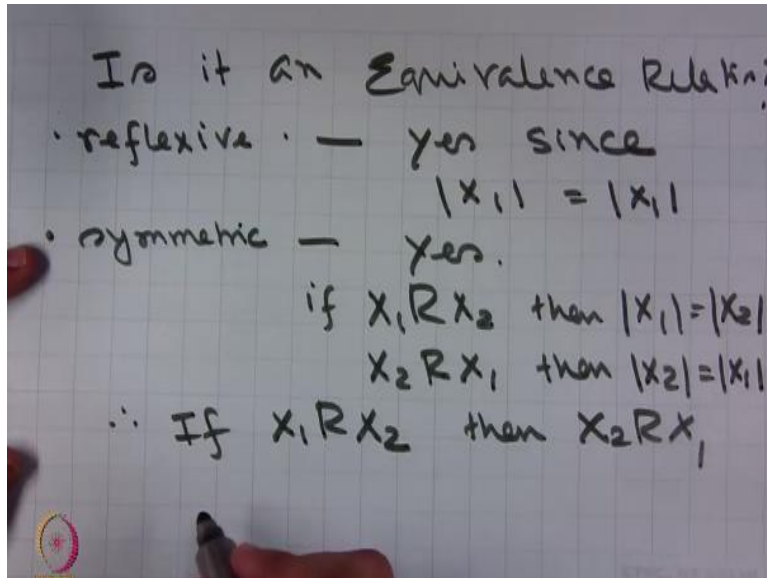
So, these are the 8 possible subsets of  $A$ .

We define  $R$  on  $X$  as follows.

For  $X_1, X_2 \in X$ ,  $X_1 R X_2$  if  $|X_1| = |X_2|$

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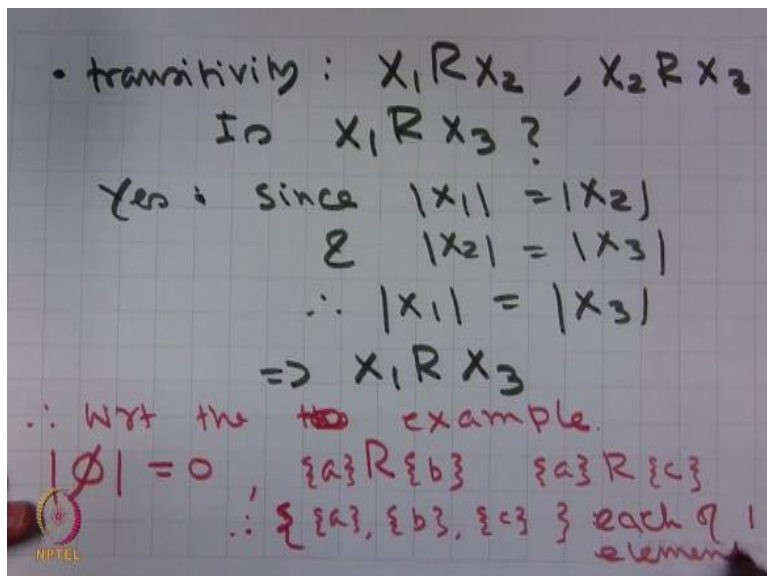


Question is, is it an Equivalence Relation?

So let us check is it

- Reflexive, yes,  
 Since,  $|x_1| = |x_1|$ . Therefore,  $x_1$  is related to itself.
- Symmetric, yes,  
 Because if  $x_1 R x_2$  then,  $|x_1| = |x_2|$  and  $x_2 R x_1$  then  $|x_2| = |x_1|$   
 Therefore, if  $x_1 R x_2$  then  $x_2 R x_1$

(Refer Slide Time: 26:27)



- Transitivity, suppose  $x_1 R x_2$  and  $x_2 R x_3$ .

Question is, is  $x_1 R x_3$ ?

Yes, since  $|x_1| = |x_2|$  and  $|x_2| = |x_3|$ . Therefore,  $|x_1| = |x_3|$

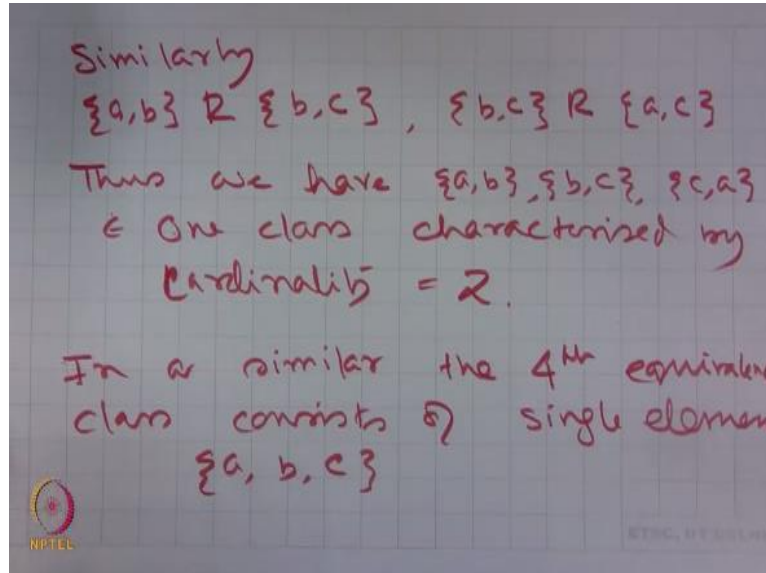
$\Rightarrow x_1 R x_3$

∴ With respect to the example we get that  $\phi$  is not related to any other set since  $|\phi| = 0$ .

$$\{a\}R\{b\}, \{a\}R\{c\}$$

∴ one class is comprising of the 3 sets  $\{a\}, \{b\}$  and  $\{c\}$ , because each has cardinality 1.

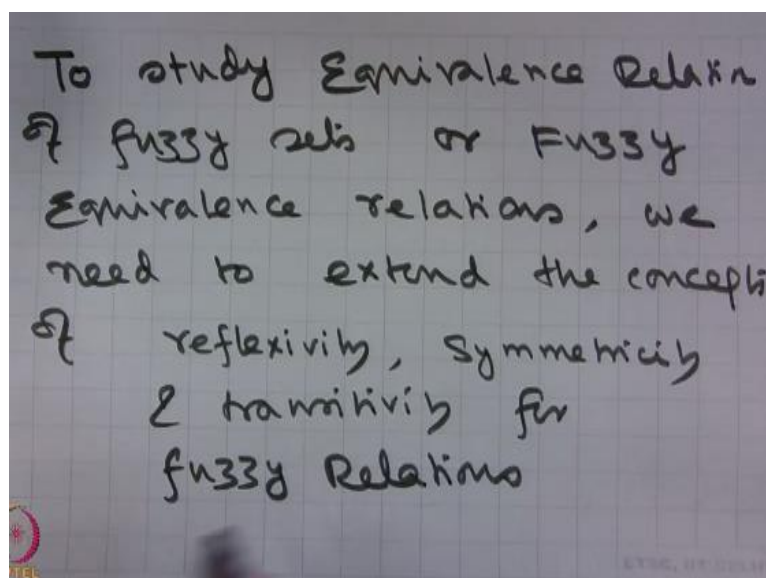
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Similarly,  $\{a,b\}R\{b,c\}$  and  $\{b,c\}R\{a,c\}$

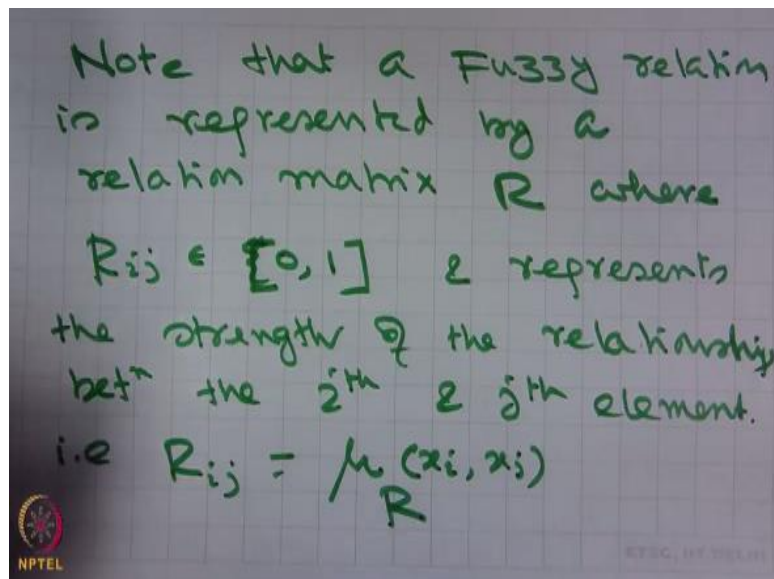
Thus, we have  $\{a,b\}, \{a,c\}$  and  $\{b,c\}$  belong to one class characterized by cardinality is equal to 2 and in a similar way the fourth equivalence class consists of single element  $\{a,b,c\}$ , which has 3 elements.

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Now, to study Equivalence Relation on fuzzy sets or Fuzzy Equivalence relations, we need to extend the concepts of the 3 properties reflexivity, symmetricity, and transitivity for fuzzy relations.

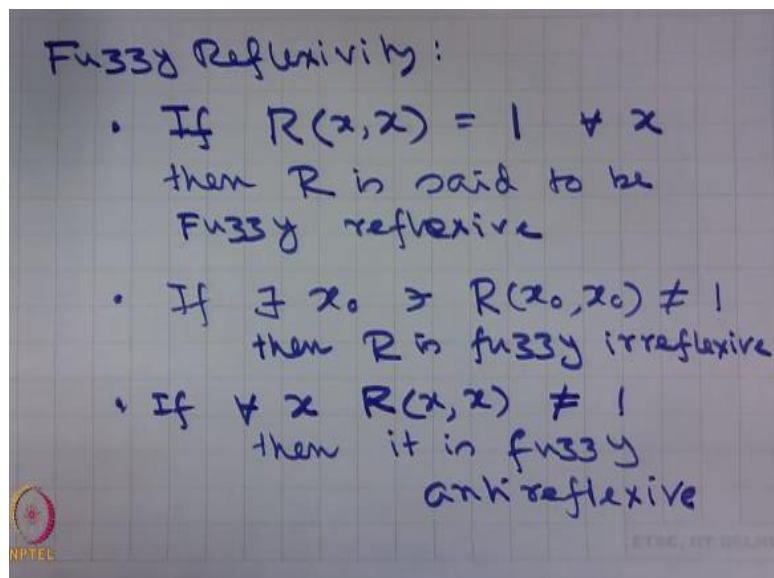
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Note that a fuzzy relation is represented by a relation matrix  $R$  where  $R_{ij} \in [0, 1]$  and represents the strength of the relationship between the  $i^{\text{th}}$  and  $j^{\text{th}}$  element that is

$$R_{ij} = \mu_R(x_i, x_j)$$

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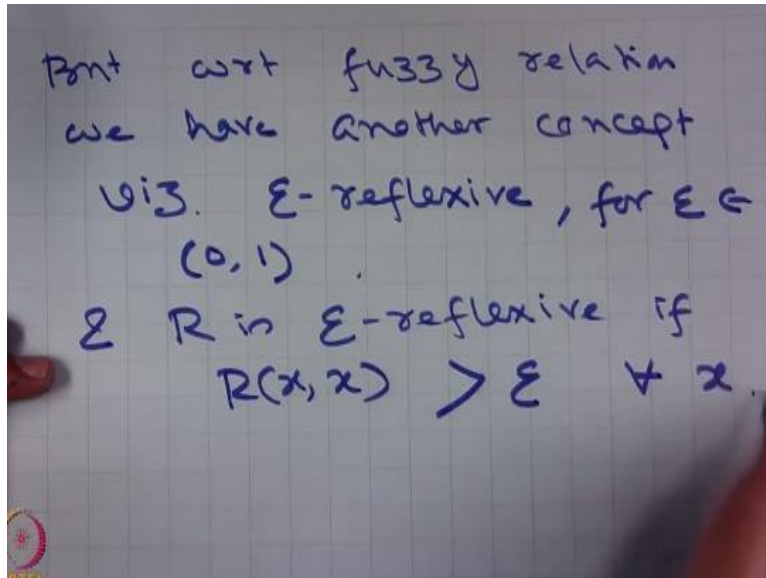


With that background let us now define Fuzzy Reflexivity.

- If  $R(x, x) = 1 \quad \forall x$  then  $R$  is said to be fuzzy reflexive.
- If  $\exists x_0$  such that  $R(x_0, x_0) \neq 1$  then  $R$  is fuzzy irreflexive.
- If  $\forall x, R(x, x) \neq 1$ , then it is fuzzy antireflexive.

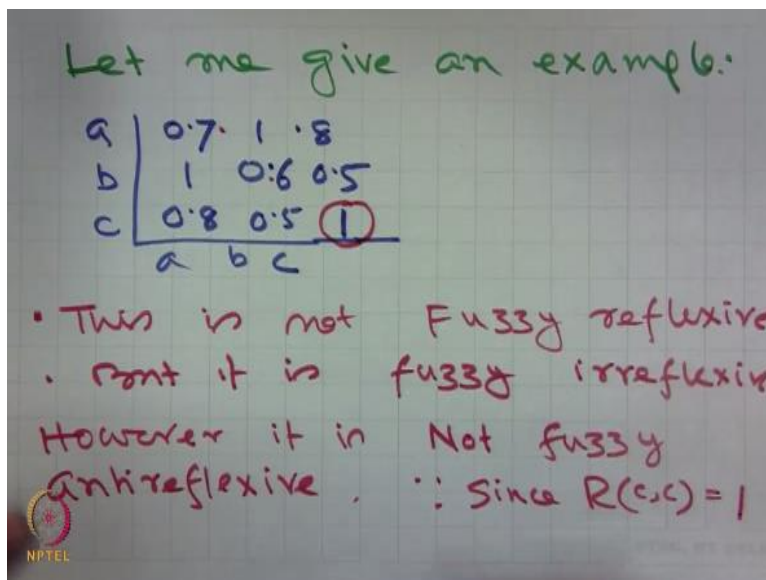
So these 3 are similar to what we define with respect to crisp relations.

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But with respect to fuzzy relation we have another concept namely  $\epsilon$ -reflexive for  $\epsilon \in [0, 1]$  and  $R$  is  $\epsilon$ -reflexive if  $R(x, x) > \epsilon \forall x$

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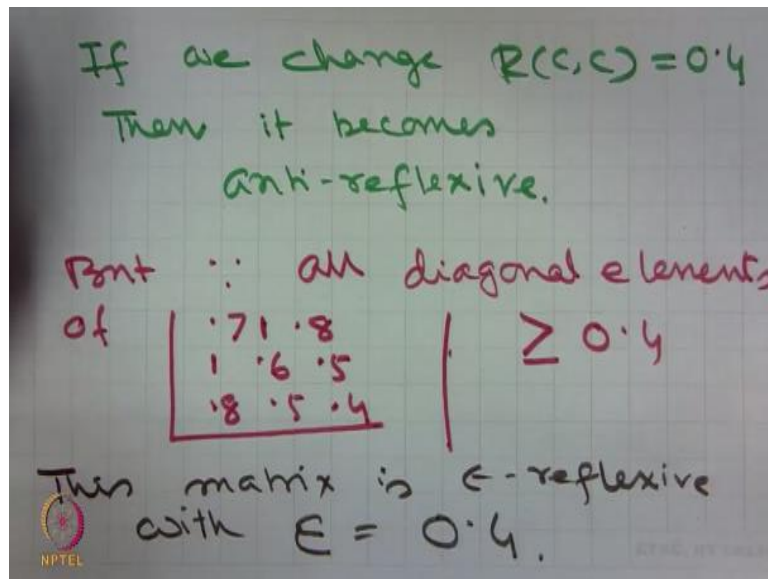


So let me give an example:

a	0.7	1	0.8
b	1	0.6	0.5
c	0.8	0.5	1
	a	b	c

- This is not fuzzy reflexive because some diagonal elements are not 1,
- But it is fuzzy irreflexive; however, it is not fuzzy anti-reflexive since  $R(c, c) = 1$

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If we change  $R(c, c) = 0.4$  then it becomes anti-reflexive.

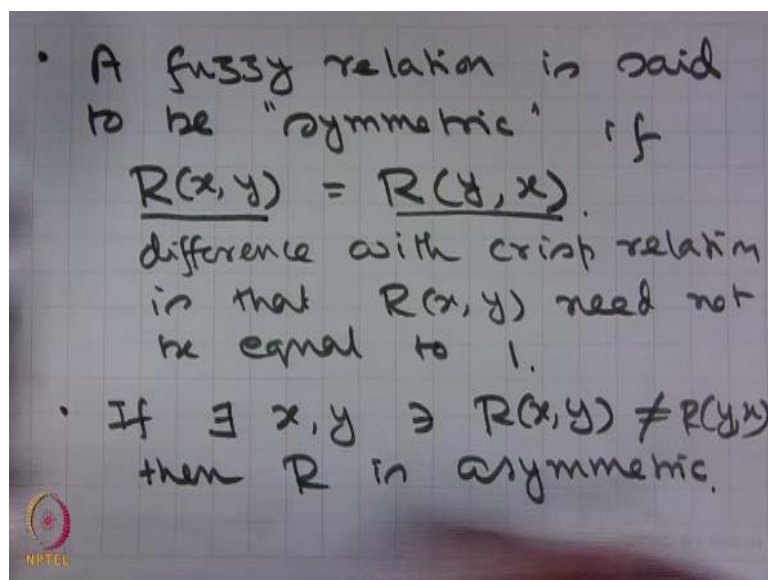
But since all diagonal elements of

$a$	0.7	1	0.8
$b$	1	0.6	0.5
$c$	0.8	0.5	0.4
	$a$	$b$	$c$

are greater than equal to 0.4.

This matrix is  $\epsilon$ -reflexive with  $\epsilon = 0.4$

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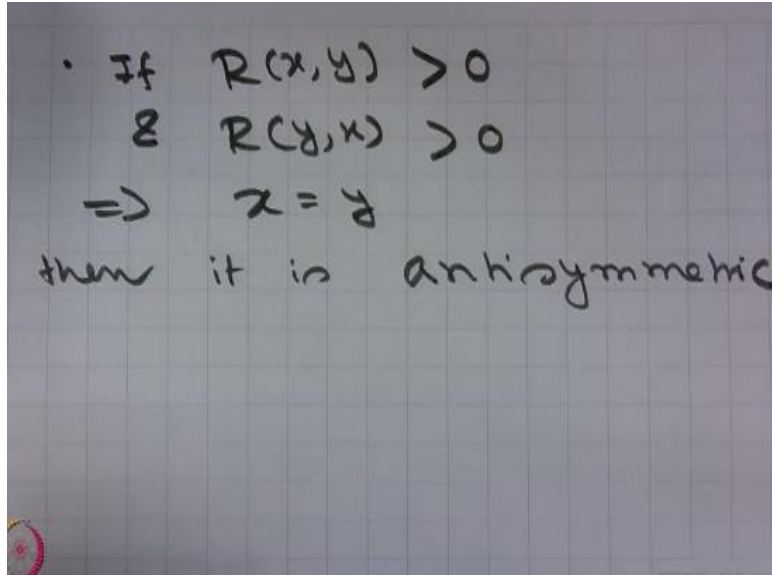
- A fuzzy relation is said to be symmetric if  $R(x,y) = R(y,x)$ . With respect to crisp relation also we have seen this, but the difference is that  $R(x,y)$  need not be equal to



1. That is if the strength of relationship between  $x$  and  $y$  is same as the strength of the relationship between  $y$  and  $x$ , whatever value it is we call it symmetric.

- Similarly, if there exists  $x, y$  such that  $R(x, y) \neq R(y, x)$  then  $R$  is asymmetric.

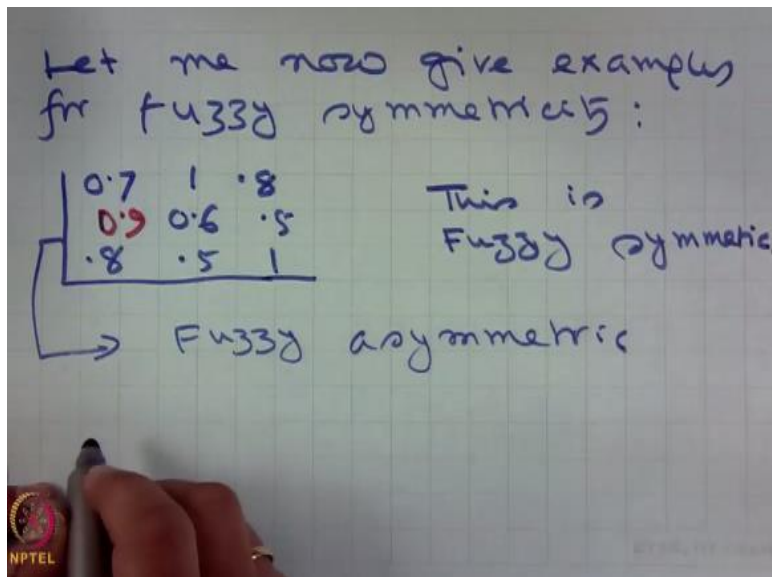
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- If  $R(x, y) > 0$  and  $R(y, x) > 0$  implies  $x = y$ , then it is anti-symmetric.

Thus we find a subtle difference with respect to anti-symmetry between crisp relation and a fuzzy relation.

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Let me now give examples for fuzzy symmetry.

Consider

$a$	0.7	1	0.8
$b$	1	0.6	0.5



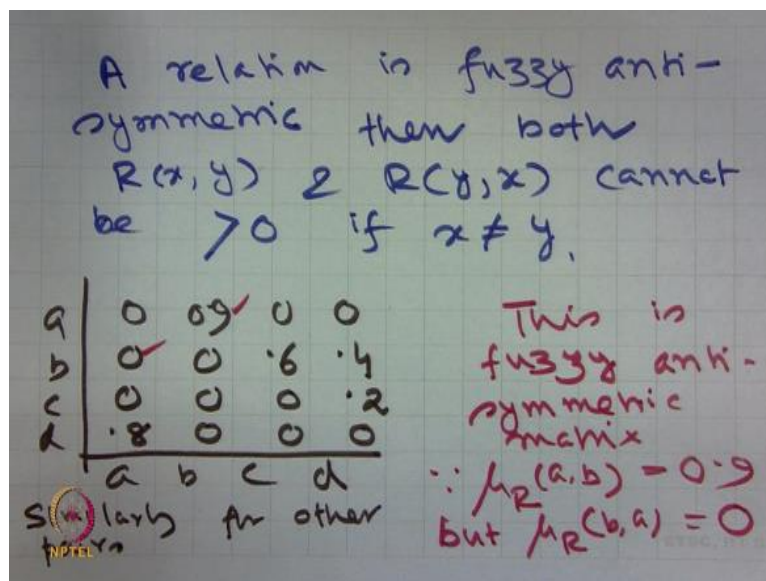
$c$	0.8	0.5	1
	$a$	$b$	$c$

This is fuzzy symmetric.

If we change  $R(b, a)$  to 0.9 then it becomes fuzzy asymmetric.

$a$	0.7	1	0.8
$b$	0.9	0.6	0.5
$c$	0.8	0.5	1
	$a$	$b$	$c$

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A relation is fuzzy anti-symmetric then both  $R(x, y)$  and  $R(y, x)$  cannot be  $> 0$ , if  $x \neq y$

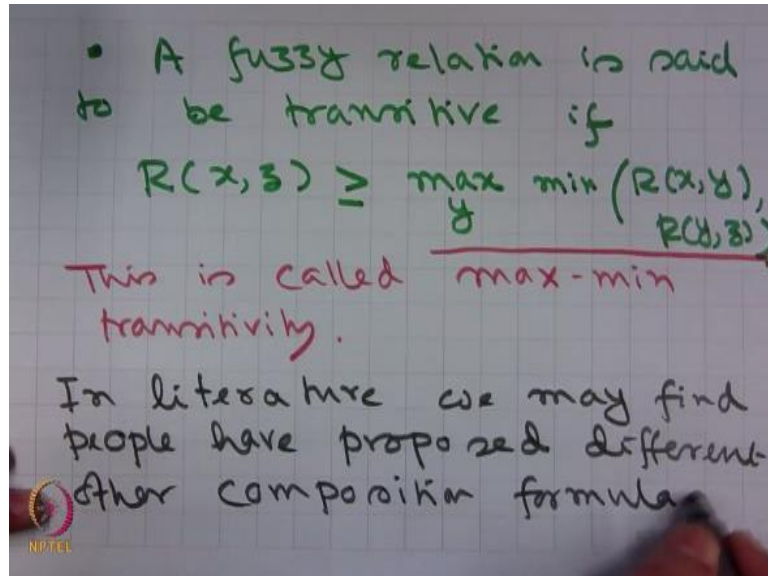
Therefore, consider this matrix

$a$	0	0.9	0	0
$b$	0	0	0.6	0.4
$c$	0	0	0	0.2
$d$	0.8	0	0	0
	$a$	$b$	$c$	$d$

This is a fuzzy anti-symmetric matrix.

Since  $\mu_R(a, b) = 0.9$  but  $\mu_R(b, a) = 0$ . Similarly, for other pairs. Please verify that it is fuzzy anti-symmetric.

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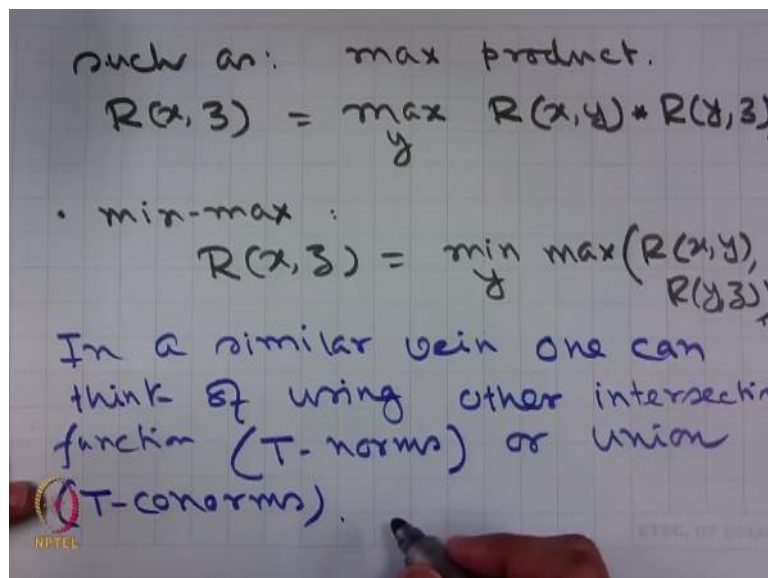
- A fuzzy relation is said to be transitive if

$$R(x, z) \geq \max_y \{\min(R(x, y), R(y, z))\}$$

So this is called max-min transitivity because in the composition of  $R(x, y)$  and  $R(y, z)$  we are using the max-min formula. Min is actually the standard intersection and max gives standard union.

In literature we may find people have proposed different other composition formula.

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Such as max-product that is

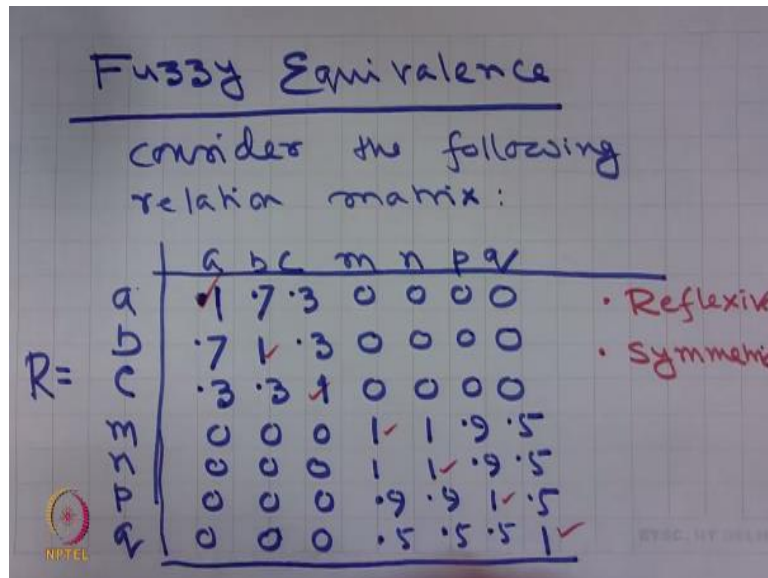
$$R(x, z) = \max_y \{R(x, y) \cdot R(y, z)\}$$

there can be min-max that is

$$R(x, z) = \max_y \{\min(R(x, y), R(y, z))\}$$

So in different applications people have tried different way of relation composition, but max min is the most commonly practiced one. In a similar vein one can think of using other intersection function that is t-norms or union that is t-conorms, but in this course we are not going into much details with respect to these different formulae.

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With that background let us now explain Fuzzy Equivalence.

Consider the following relation matrix.

	a	b	c	m	n	p	q
a	1	0.7	0.3	0	0	0	0
b	0.7	1	0.3	0	0	0	0
c	0.3	0.3	1	0	0	0	0
m	0	0	0	1	1	0.9	0.5
n	0	0	0	1	1	0.9	0.5
p	0	0	0	0.9	0.9	1	0.5
q	0	0	0	0.5	0.5	0.5	1

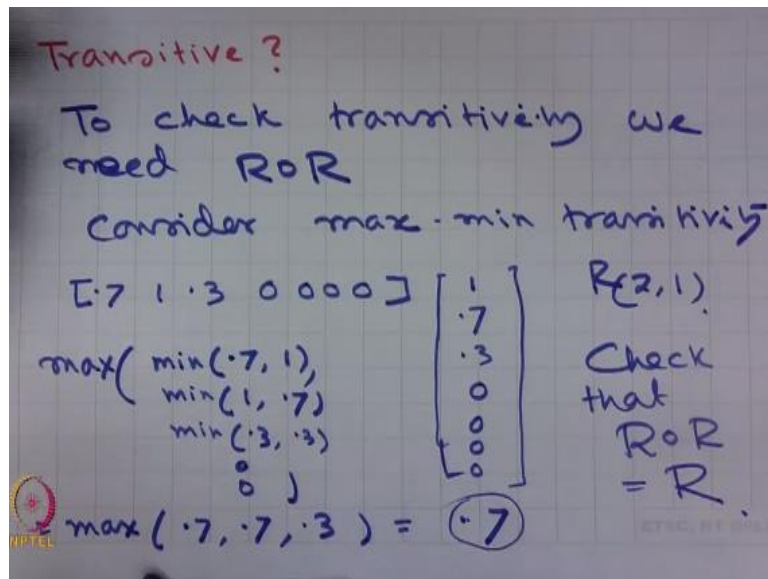
So let us call it R.

If we examine the matrix carefully, we see that all the diagonal elements are 1. Therefore, it is reflexive. If you recall the definition that a fuzzy relation is reflexive if  $R(x, x) = 1$  for all x.

Also note that it is symmetric.

Now is it transitive?

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To check transitivity, we need to compose  $R$  with itself that is,  $R \circ R$  and suppose we work on max-min transitivity.

Let us consider any arbitrary element say  $ij$  and if we compose what we shall get. So let us look at the say second row first column element.

The second row is  $[0.7 \ 1 \ 0.3 \ 0 \ 0 \ 0 \ 0]$  and the first column is  $\begin{bmatrix} 1 \\ 0.7 \\ 0.3 \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$ . What is the max-min?

$$R \circ R(2,1) = \max \left\{ \begin{array}{l} \min(0.7,1) \\ \min(1,0.7) \\ \min(0.3,0.3) \\ \min(0,0) \\ \min(0,0) \\ \min(0,0) \\ \min(0,0) \end{array} \right\} = \max\{0.7,0.3,0\} = 0.7$$

In the original matrix the 2, 1 element was 0.7 and after composition also it remains 0.7.

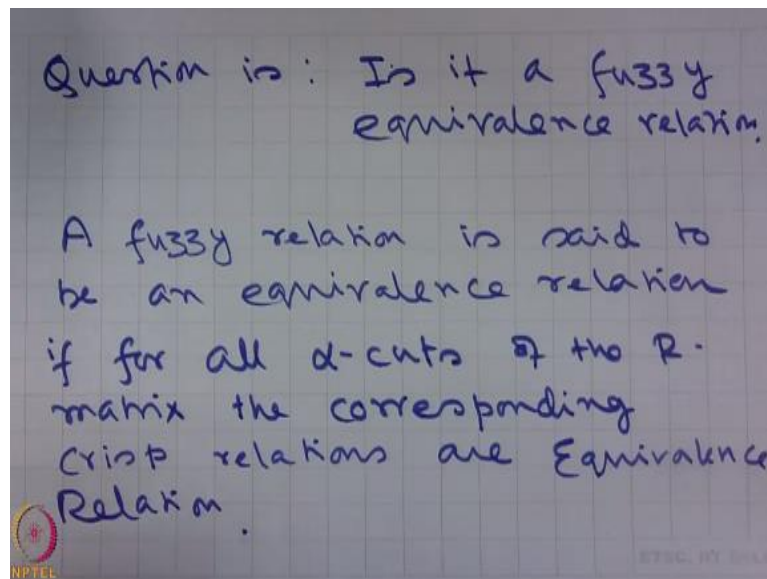
I request you to verify for all the 49 elements and see that  $R \circ R = R$

That is the transitive closure of  $R$  if we apply that algorithm that I have given some time back after the first iteration itself we find that  $R'$  is same as  $R$  and therefore, this relation is transitive.

Also if we look at the matrix very carefully you can see that the first 3 elements are related with each other and the last 4 elements are related with each other, these have been kept 0.

Therefore, clear distinction between the elements which are related with each other.

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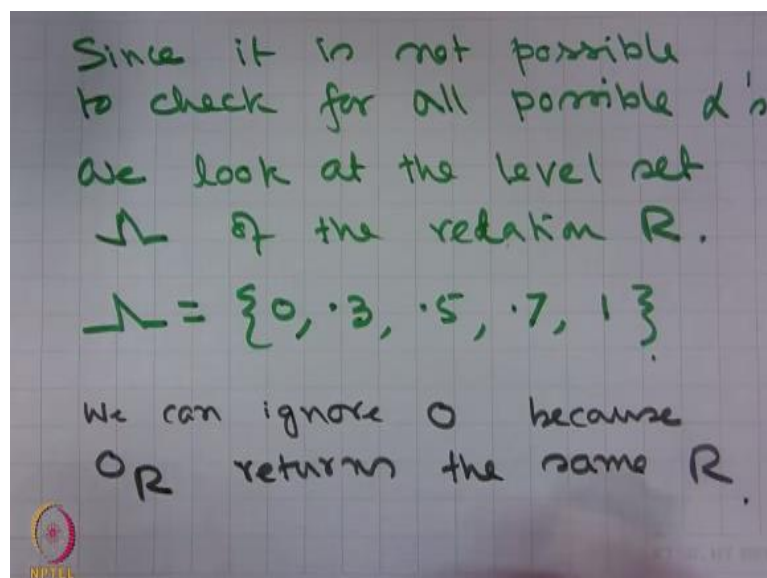


Therefore, the question is, is it a fuzzy equivalence relation?

A fuzzy relation is said to be an equivalence relation if for all  $\alpha$ -cuts of the  $R$  matrix, the corresponding crisp relations are equivalence relation.

Therefore, if we have been given a fuzzy relation matrix then corresponding to all  $\alpha$ -cuts we need to verify if we are getting an equivalence relation among the elements of the underlying set.

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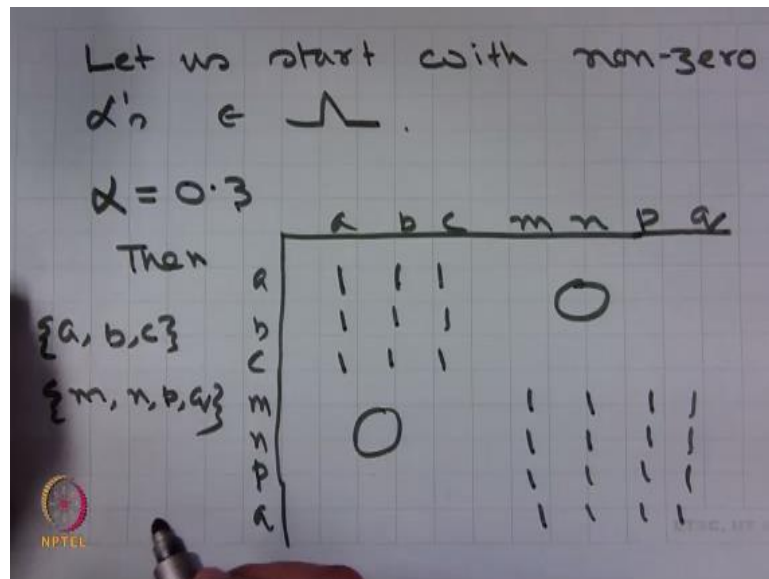


Since it is not possible to check for all possible  $\alpha$ 's.

We look at the level set ( $\Lambda$ ) of the relation  $R$ . Now if we look at the matrix very carefully then we get that different  $\alpha$ 's belonging to the level set  $\Lambda = \{0, 0.3, 0.5, 0.7, 1\}$ . Therefore, we need

to check for different alphas and see what happens. We can ignore 0, because  ${}^0R$  returns the same  $R$ .

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So let us start with non-zero  $\alpha$ 's belonging to the level set  $\Lambda$ .

Consider  $\alpha = 0.3$  then, the matrix that we will get is

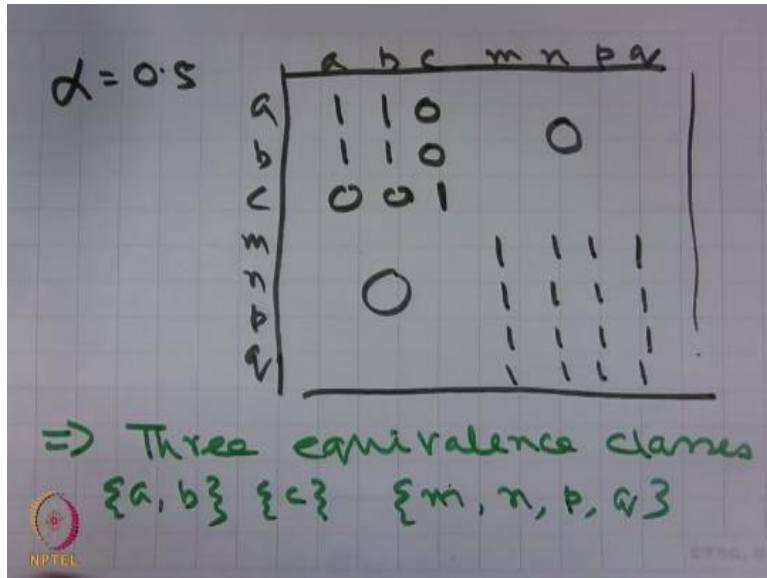
	$a$	$b$	$c$	$m$	$n$	$p$	$q$
$a$	1	1	1	0	0	0	0
$b$	1	1	1	0	0	0	0
$c$	1	1	1	0	0	0	0
$m$	0	0	0	1	1	1	1
$n$	0	0	0	1	1	1	1
$p$	0	0	0	1	1	1	1
$q$	0	0	0	1	1	1	1

It is very clear that these are 2 different equivalence classes.

Therefore, we get that 2 classes to be  $\{a, b, c\}$  and  $\{m, n, p, q\}$ .

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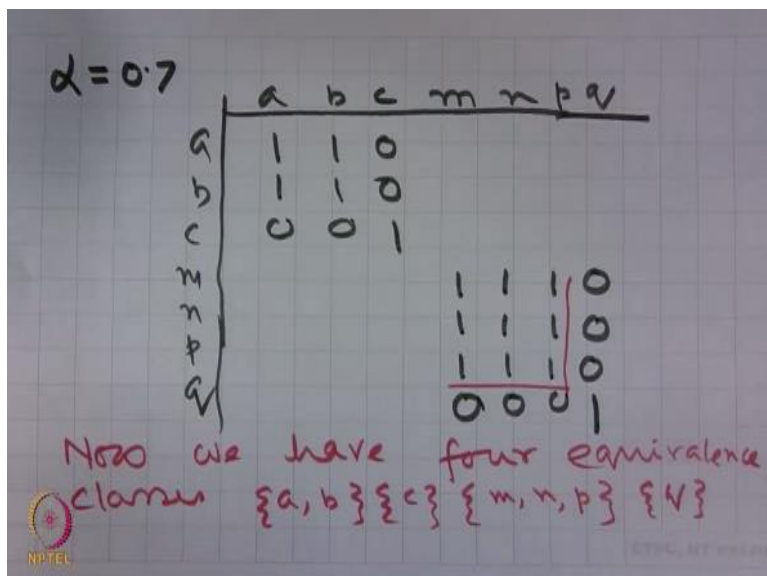


Let us now consider  $\alpha = 0.5$ , therefore, the matrix that we get is,

	a	b	c	m	n	p	q
a	1	1	0	0	0	0	0
b	1	1	0	0	0	0	0
c	0	0	1	0	0	0	0
m	0	0	0	1	1	1	1
n	0	0	0	1	1	1	1
p	0	0	0	1	1	1	1
q	0	0	0	1	1	1	1

So that gives us 3 equivalence classes namely  $\{a, b\}$ ,  $\{c\}$  and  $\{m, n, p, q\}$ .

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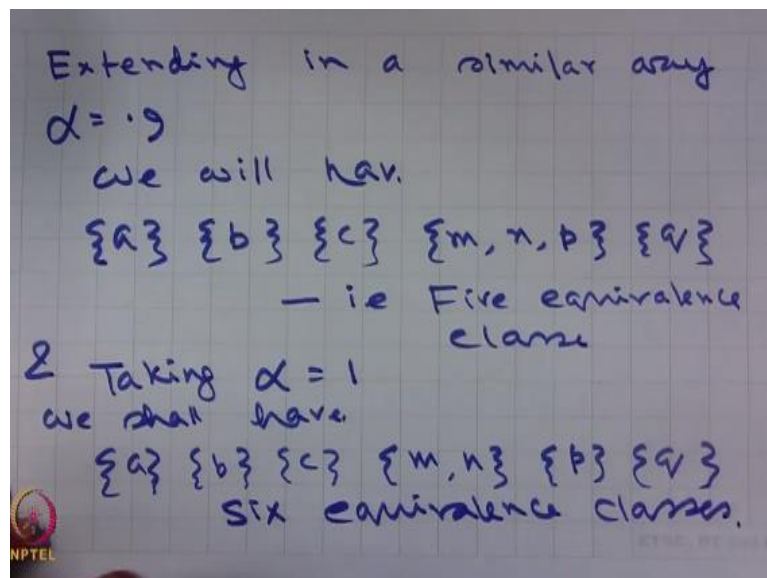


When we take  $\alpha = 0.7$ ,

	$a$	$b$	$c$	$m$	$n$	$p$	$q$
$a$	1	1	0	0	0	0	0
$b$	1	1	0	0	0	0	0
$c$	0	0	1	0	0	0	0
$m$	0	0	0	1	1	1	0
$n$	0	0	0	1	1	1	0
$p$	0	0	0	1	1	1	0
$q$	0	0	0	0	0	0	1

Thus, with  $\alpha = 0.7$ , now we have 4 equivalence classes namely  $\{a, b\}$ ,  $\{c\}$ ,  $\{m, n, p\}$  and  $\{q\}$ .

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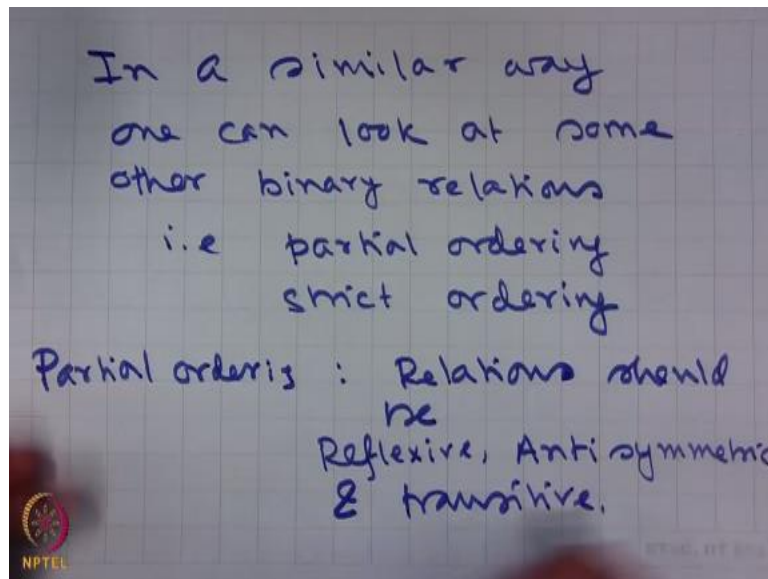


Extending in a similar way for  $\alpha = 0.9$ , we will have  $\{a\}$ ,  $\{b\}$ ,  $\{c\}$ ,  $\{m, n, p\}$  and  $\{q\}$  that is, 5 equivalence classes and taking  $\alpha = 1$  we shall have  $\{a\}$ ,  $\{b\}$ ,  $\{c\}$ ,  $\{m, n\}$ ,  $\{p\}$  and  $\{q\}$  that is 6 equivalence classes.

Thus what we find that for different  $\alpha$ -cuts we get the corresponding partitions and each one of them represent equivalence classes.

Therefore, we can say that this fuzzy relation is actually a fuzzy equivalence relation.

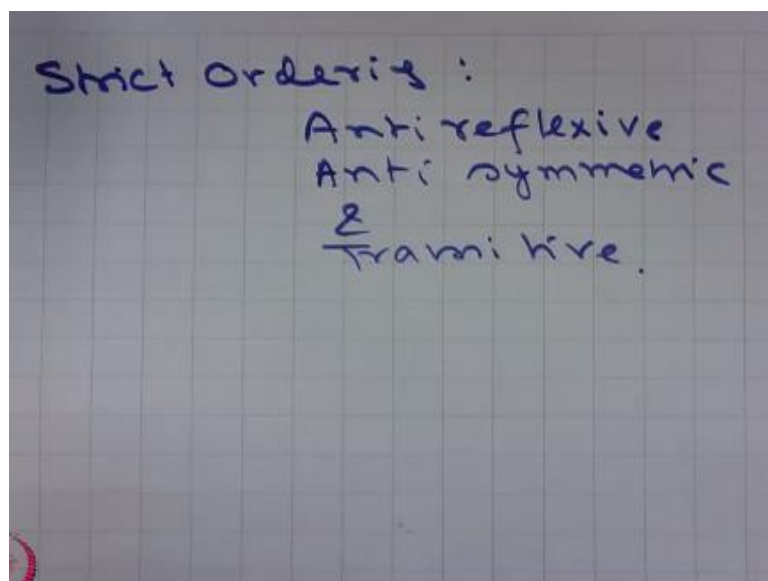
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In a similar way one can look at some other binary relations namely partial ordering, strict ordering.

For partial ordering we have seen the relation should be reflexive anti-symmetric and transitive.

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For strict ordering that it should be anti-reflexive, anti-symmetric and transitive.

We have already extended the definitions of these properties with respect to fuzzy relations.

So one can apply them for a given relation matrix and examine their properties. Okay students

I stop here today and with this I complete my chapter on relations, from the next class onwards

I shall start fuzzy logic which is an important concept with respect to decision making and

artificial intelligence. Thank you.